

The Omega-Reducibility of Certain Pseudovarieties of Ordered Monoids

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- 5 Briefly about the ω -reducibility of more complex pseudovarieties

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there exists an inequality $u' \leq v'$ of ω -words that is also valid in V and “has the same α -imprint in M ”, i.e.,

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Construct the V_0 -completion α_{V_0} of the homomorphism α :

- $\alpha_{V_0}: \widehat{A}^* \rightarrow M \times A^* / \equiv_{V_0}$
- $x \mapsto (\alpha(x), [x]_{\equiv_{V_0}})$

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- Proof of the ω -reducibility of $V_{1/2}$ inspired by the approach of Place and Zeitoun in their papers (2014–2019) on certain properties of levels of concatenation hierarchies.

GIVEN:

- locally finite pseudovariety V_0 ,
- finite alphabet A ,
- finite ordered monoid M ,
- continuous homomorphism $\alpha: \widehat{A}^* \rightarrow M$.

Construct the V_0 -completion α_{V_0} of the homomorphism α :

- $\alpha_{V_0}: \widehat{A}^* \rightarrow M \times A^* / \equiv_{V_0}$
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Construct the image $M_{\alpha_{V_0}}$ of α_{V_0} .

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OUTPUT of Step 1:

V_0 -compatible onto continuous homomorphism $\alpha_{V_0}: \widehat{A}^* \rightarrow M_{\alpha_{V_0}}$:

$\forall x, y \in \widehat{A}^*: \alpha_{V_0}(x) = \alpha_{V_0}(y) \Rightarrow x \equiv_{V_0} y.$

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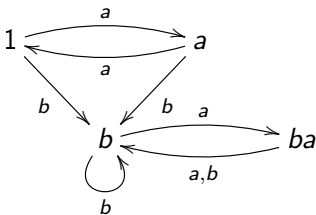
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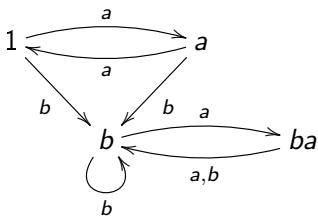
Cayley graph of monoid M :



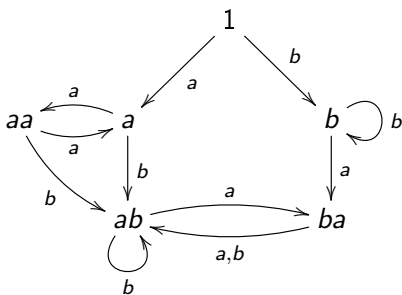
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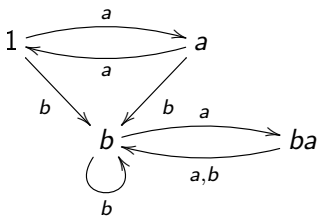
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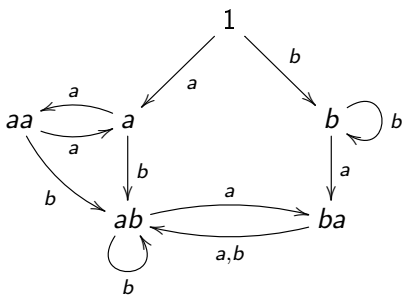
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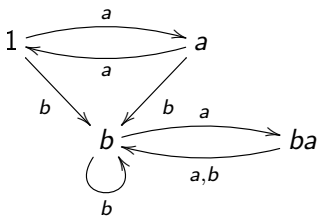
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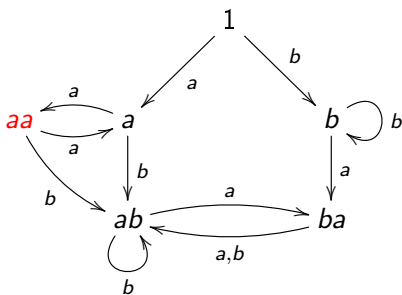
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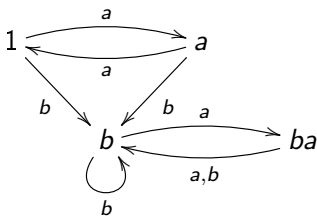
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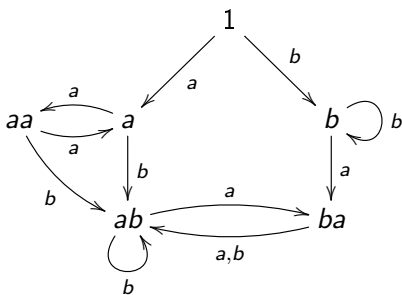
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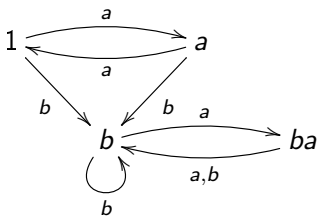
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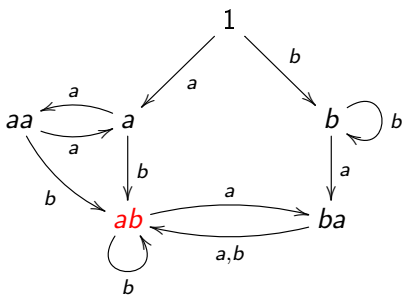
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Let $u, v \in A^*$ be arbitrary words and $N = 1 + 9 \cdot |M_{\alpha_{V_0}}| \cdot 2^{|M_{\alpha_{V_0}}|}$.
Let $V_{1/2}^N \models u \leq v$. Then there exist ω -words u', v' satisfying

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- Proven by the induction on the height of a *factorization tree* of word u for α_{V_0} .

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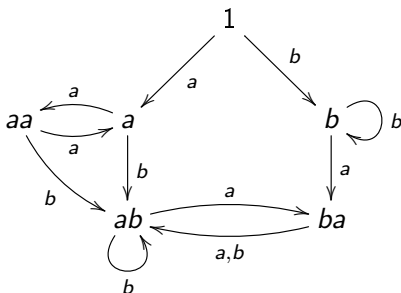
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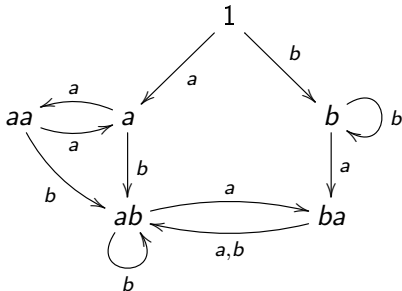


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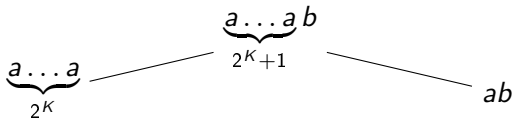
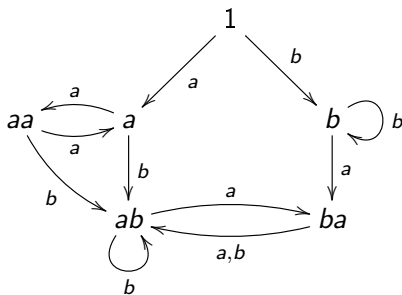
$$2^{K+1}$$

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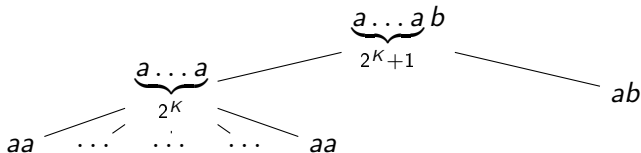
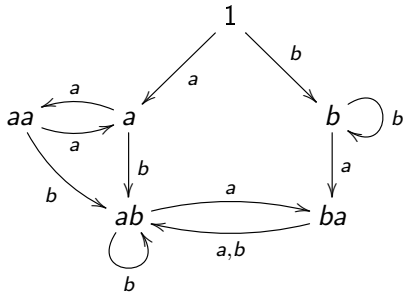


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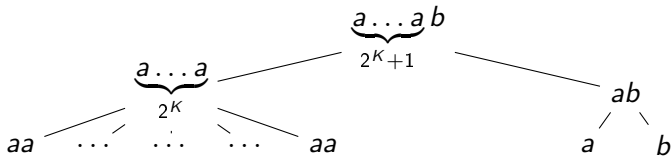
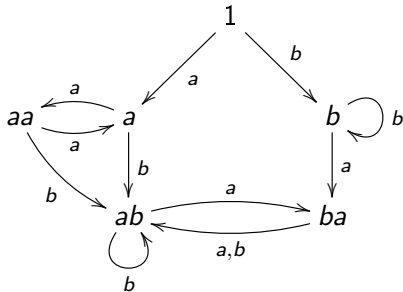


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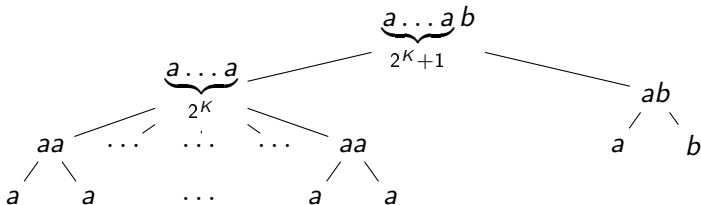
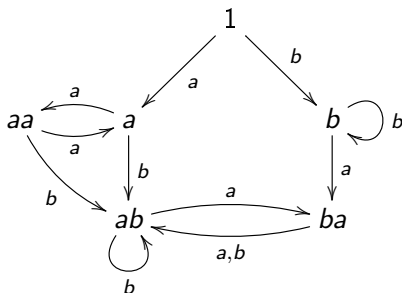


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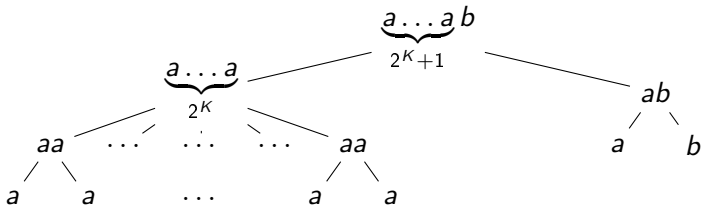
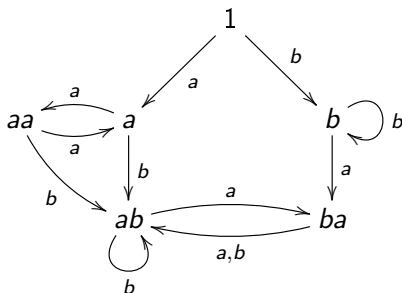


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height of this factorization tree = 4

Theorem (Simon, 1990, Kufleitner, 2008)

Let A be a finite alphabet, M a finite monoid, $\alpha: A^ \rightarrow M$ a homomorphism. Then, for every word $u \in A^*$, there exists a factorization tree of u for α of height at most $3 \cdot |M|$.*

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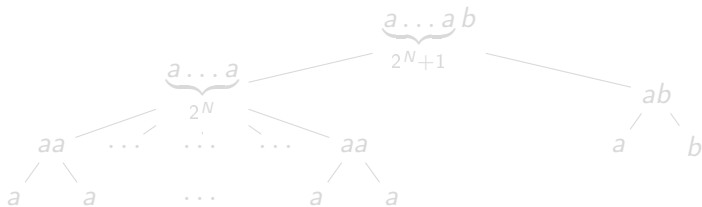
The following lemma will be useful.

Lemma (J. V.)

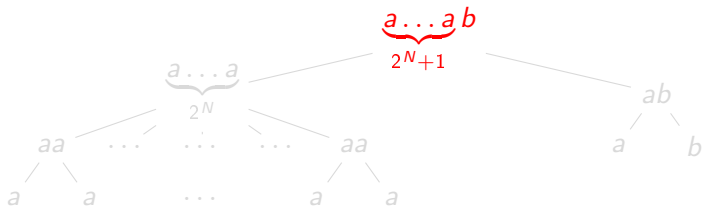
Let $u_1, \dots, u_k, v \in A^*$ ($k \in \mathbb{N}$), $n \geq k - 1$. Let $V_{1/2}^n \models u_1 \dots u_k \leq v$. Then there exist $v_1, \dots, v_k \in A^*$ such that

- $v = v_1 \dots v_k$,
- $V_{1/2}^{n-(k-1)} \models u_i \leq v_i$ for $i = 1, \dots, k$.

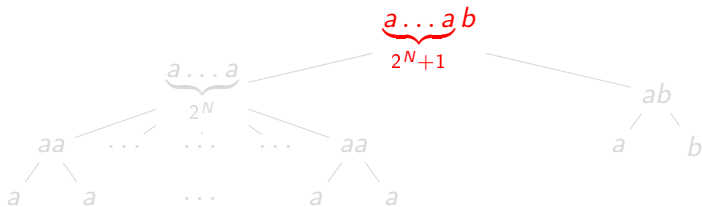
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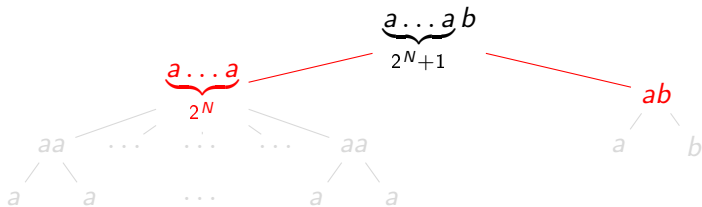


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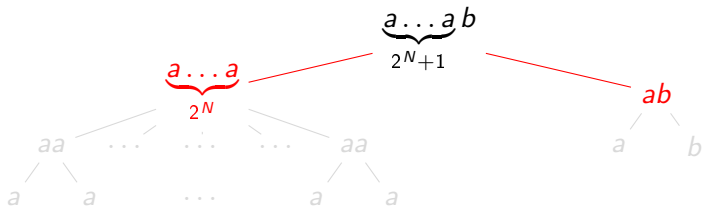
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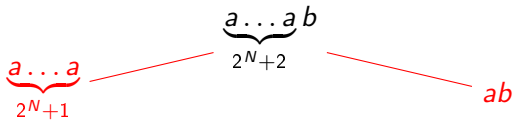
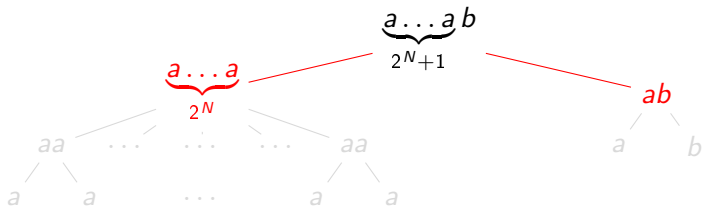
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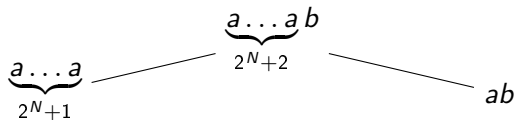
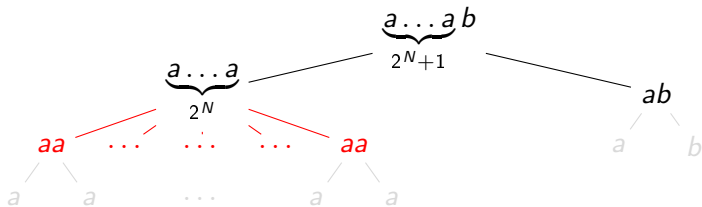


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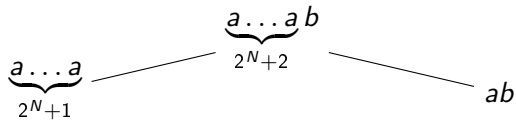
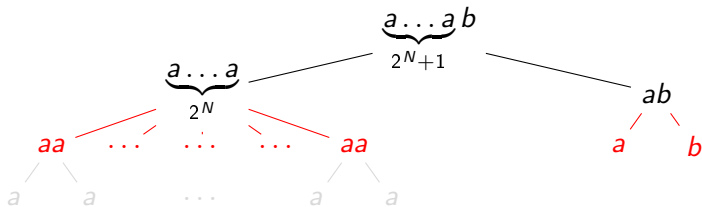
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$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b \rightarrow V_{1/2}^{N-1} \models \underbrace{a \dots a}_{2^N} \leq \underbrace{a \dots a}_{2^{N+1}}, ab \leq ab$$

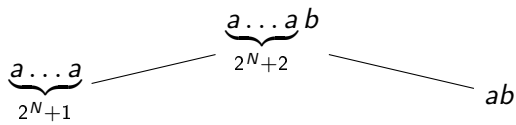
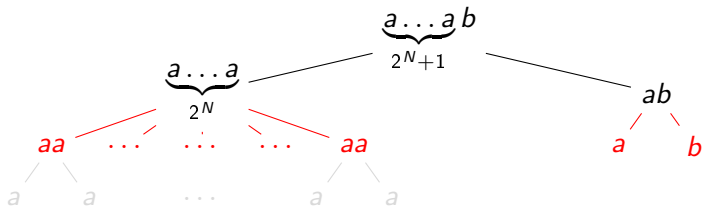


$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b \rightarrow V_{1/2}^{N-1} \models \underbrace{a \dots a}_{2^N} \leq \underbrace{a \dots a}_{2^{N+1}}, ab \leq ab$$



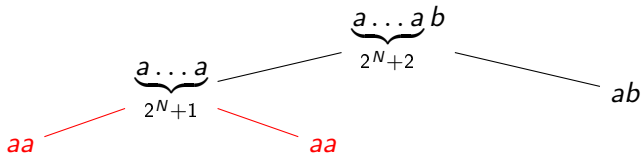
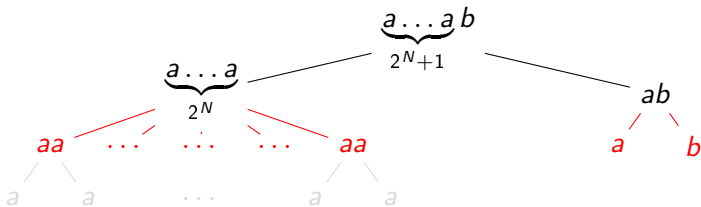
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$$\rightarrow V_{1/2}^{N-3} \models aa \leq aa, \underbrace{a \dots a}_{2^{N-4}} \leq \underbrace{a \dots a}_{2^{N-3}}, aa \leq aa, a \leq a, b \leq b$$



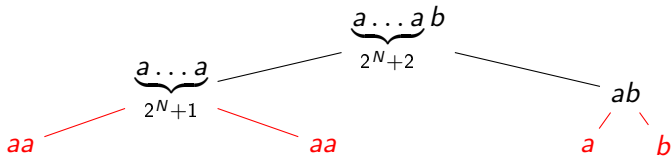
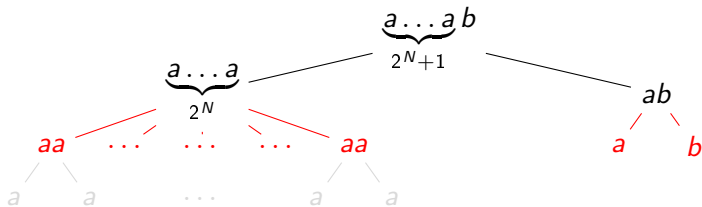
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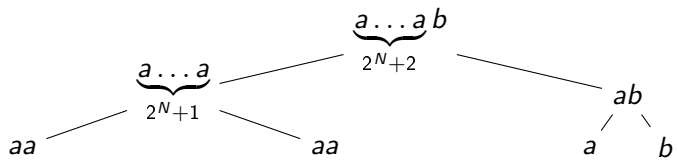
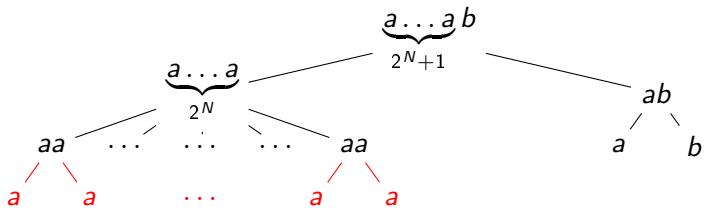
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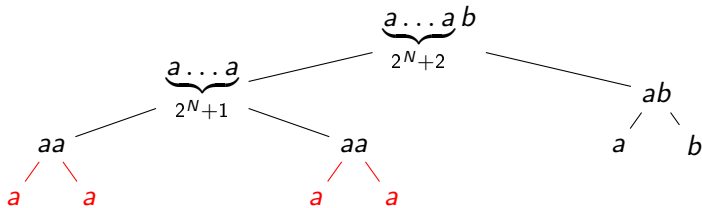
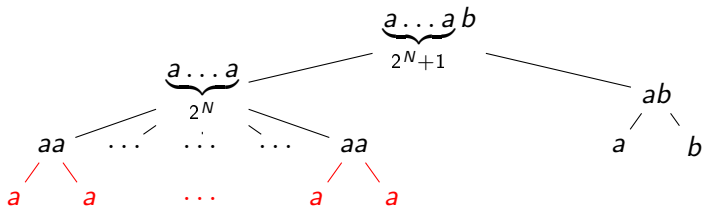
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$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b \rightarrow V_{1/2}^{N-1} \models \underbrace{a \dots a}_{2^N} \leq \underbrace{a \dots a}_{2^{N+1}}, ab \leq ab$$

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$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b$$

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$$\rightarrow V_{1/2}^{N-3} \models aa \leq aa, \underbrace{a \dots a}_{2^{N-4}} \leq \underbrace{a \dots a}_{2^{N-3}}, aa \leq aa, \quad a \leq a, b \leq b$$

$$\rightarrow V_{1/2}^{N-4} \models a \leq a, a \leq a$$

$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b$$

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$$\rightarrow V_{1/2}^{N-3} \models aa \leq aa, \underbrace{a \dots a}_{2^{N-4}} \leq \underbrace{a \dots a}_{2^{N-3}}, aa \leq aa, \quad a \leq a, b \leq b$$

$$\rightarrow V_{1/2}^{N-4} \models a \leq a, a \leq a$$

- $V_{1/2} \models a \leq a$

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- $V_{1/2} \models aa \leq aa$

$$\rightarrow V_{1/2}^{N-4} \models a \leq a, a \leq a$$

- $V_{1/2} \models a \leq a$

$$V_{1/2}^N \models \underbrace{a \dots a}_{2^{N+1}} b \leq \underbrace{a \dots a}_{2^{N+2}} b$$

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$$V_{1/2} := \llbracket u^{\omega+1} \leq u^\omega v u^\omega \mid V_0 \models u = v \rrbracket$$

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- $V_0 = \text{SI} \models aa = a$

$$\rightarrow V_{1/2}^{N-3} \models aa \leq aa, \underbrace{a \dots a}_{2^{N-4}} \leq \underbrace{a \dots a}_{2^{N-3}}, aa \leq aa, \quad a \leq a, b \leq b$$

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$$\rightarrow V_{1/2}^{N-3} \models aa \leq aa, \underbrace{a \dots a}_{2^{N-4}} \leq \underbrace{a \dots a}_{2^{N-3}}, aa \leq aa, \quad a \leq a, b \leq b$$

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ω -Reducibility of $V_{3/2}$ and $V_{5/2}$

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$V_{1/2}$
 ω -reducibility
of *pairs* of words

v
| 1/2
u

ω -Reducibility of $V_{3/2}$ and $V_{5/2}$

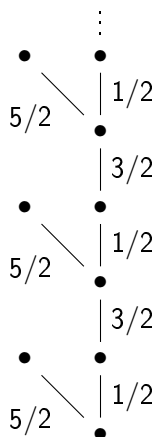
$V_{1/2}$
 ω -reducibility
of *pairs* of words

v
| 1/2
u

$V_{3/2}$
 ω -reducibility
of *chains* of words

w
| 3/2
v
| 1/2
u

ω -Reducibility of $V_{3/2}$ and $V_{5/2}$

$V_{1/2}$ ω -reducibility of <i>pairs</i> of words	$V_{3/2}$ ω -reducibility of <i>chains</i> of words	$V_{5/2}$ ω -reducibility of (finite) <i>ordered</i> sets of words
$\begin{array}{c} v \\ \\ u \end{array} \quad 1/2$	$\begin{array}{c} w \\ \\ v \\ \\ u \end{array} \quad \begin{array}{c} 3/2 \\ 1/2 \end{array}$	 <p>The diagram illustrates the structure of $V_{5/2}$ as a sequence of nodes. A vertical line of nodes is shown, with vertical segments between them labeled $1/2, 3/2, 1/2, 3/2, 1/2$ from top to bottom. Diagonal lines connect nodes in a zig-zag pattern, with segments labeled $5/2$. Vertical ellipsis dots at the top and bottom indicate the sequence continues infinitely in both directions.</p>

Thank you for your attention.