

# Abstract commutator theory in concrete classes

David Stanovský

Charles University, Prague, Czech Republic

Summer School on Algebras and Ordered Sets, September 2023

Part I:

Solvability and nilpotence: the beginnings and motivation

Part II:

Abelianness and Centrality: examples and module representation

Part III:  
The commutator in specific  
varieties

# The commutator

*Centralizing relation* for congruences  $\alpha, \beta, \delta$  of an algebra  $A$ :

$C(\alpha, \beta; \delta)$  iff

for every term  $t(x_1, \dots, x_m, y_1, \dots, y_n)$  and every  $a \stackrel{\alpha}{\equiv} b$ ,  $u_i \stackrel{\beta}{\equiv} v_i$

$$t(a, \bar{u}) \stackrel{\delta}{\equiv} t(a, \bar{v}) \Rightarrow t(b, \bar{u}) \stackrel{\delta}{\equiv} t(b, \bar{v})$$

The *commutator*  $[\alpha, \beta]$  is the smallest  $\delta$  such that  $C(\alpha, \beta; \delta)$ .

# The commutator

*Centralizing relation* for congruences  $\alpha, \beta, \delta$  of an algebra  $A$ :

$C(\alpha, \beta; \delta)$  iff

for every term  $t(x_1, \dots, x_m, y_1, \dots, y_n)$  and every  $a \equiv_{\alpha} b$ ,  $u_i \equiv_{\beta} v_i$

$$t(a, \bar{u}) \equiv_{\delta} t(a, \bar{v}) \Rightarrow t(b, \bar{u}) \equiv_{\delta} t(b, \bar{v})$$

The *commutator*  $[\alpha, \beta]$  is the smallest  $\delta$  such that  $C(\alpha, \beta; \delta)$ .

A congruence  $\alpha$  of  $A$  is called

- *abelian* if  $C(\alpha, \alpha; 0_A)$ , i.e., if  $[\alpha, \alpha] = 0_A$ .
- *central* if  $C(\alpha, 1_A; 0_A)$ , i.e., if  $[\alpha, 1_A] = 0_A$ .

An algebra  $A$  is called *solvable*, resp. *nilpotent*, if there are congruences  $\alpha_i$  such that

$$0_A = \alpha_0 \leq \alpha_1 \leq \dots \leq \alpha_k = 1_A$$

and  $\alpha_{i+1}/\alpha_i$  is an *abelian*, resp. *central* congruence of  $A/\alpha_i$ , for all  $i$ .

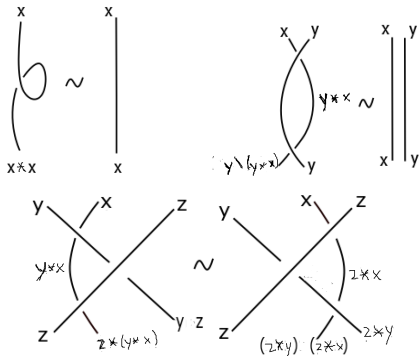
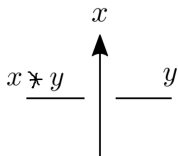
# QUANDLES

# Quandles

A binary algebra  $(Q, *, \setminus)$  is called a *quandle* if

- (I)  $x * x = x$
- (II) for all  $x, y$  there is a unique  $u$  such that  $x * u = y$
- (III)  $x * (y * z) = (x * y) * (x * z)$

Motivation: knot theory



# Quandles

A binary algebra  $(Q, *, \setminus)$  is called a *quandle* if

- (I)  $x * x = x$
- (II) for all  $x, y$  there is a unique  $u$  such that  $x * u = y$
- (III)  $x * (y * z) = (x * y) * (x * z)$

## Examples:

- group conjugation:  $(G, *)$ ,  $x * y = xyx^{-1}$
- affine forms over abelian groups:  $(G, *)$ ,  $x * y = (1 - f)(x) + f(y)$
- ...
- latin quandles = (left) self-distributive quasigroups [since 1923!]

## Quandles and associated groups

A binary algebra  $(Q, *, \setminus)$  is called a *quandle* if

- $x * x = x$
- all left translations  $L_x(y) = x * y$  are automorphisms.

*Left multiplication group, displacement group:*

$$\text{LMlt}(Q) = \langle L_x : x \in Q \rangle \leq \text{Aut}(Q)$$

$$\text{Dis}(Q) = \langle L_x L_y^{-1} : x, y \in Q \rangle \leq \text{LMlt}(Q)$$

$Q$  is *connected* if  $\text{LMlt}(Q)$  acts transitively.



# Commutator theory for quandles

[Bonatto, S. 2021]

Let  $N(Q) = \{N \leq \text{Dis}(Q) : N \text{ is normal in } \text{LMlt}(Q)\}$

There is a Galois correspondence

$$\begin{aligned} \text{Con}(Q) &\longleftrightarrow N(Q) \\ \alpha &\rightarrow \text{Dis}_\alpha = \langle L_x L_y^{-1} : x \alpha y \rangle \\ \alpha_N &= \{(x, y) : L_x L_y^{-1} \in N\} \leftarrow N \end{aligned}$$

# Commutator theory for quandles

[Bonatto, S. 2021]

Let  $N(Q) = \{N \leq \text{Dis}(Q) : N \text{ is normal in } \text{LMlt}(Q)\}$

There is a Galois correspondence

$$\begin{aligned} \text{Con}(Q) &\longleftrightarrow N(Q) \\ \alpha &\rightarrow \text{Dis}_\alpha = \langle L_x L_y^{-1} : x \alpha y \rangle \\ \alpha_N &= \{(x, y) : L_x L_y^{-1} \in N\} \leftarrow N \end{aligned}$$

## Proposition

*TFAE for congruences  $\alpha, \beta$  of a quandle  $Q$ :*

- 1  $\alpha$  centralizes  $\beta$  over  $0_Q$ , i.e.,  $C(\alpha, \beta; 0_Q)$
- 2  $[\text{Dis}_\alpha, \text{Dis}_\beta] = 1$  and  $\text{Dis}_\beta$  acts  $\alpha$ -semiregularly on  $Q$

$\alpha$ -semiregularly means  $g(a) = a \Rightarrow g(b) = b$  for every  $b \stackrel{\alpha}{\equiv} a$

# Abelian, nilpotent, and solvable quandles

[Jedlička, Pilitowska, S, Zamojska-Dzienio, 2018] [Bonatto, S, 2021]

quandle		$\text{Dis}(Q)$
affine	$\Leftrightarrow$	abelian, semiregular, "balanced"
$\Downarrow$		$\Downarrow$
abelian	$\Leftrightarrow$	abelian, semiregular
$\Downarrow$		$\Downarrow$
nilpotent	$\Leftrightarrow$	nilpotent
$\Downarrow$		$\Downarrow$
solvable	$\Leftrightarrow$	solvable

Moreover, for **finite connected faithful** quandles:

nilpotent  $\Rightarrow$  direct product of connected quandles of **prime power size**.

## Application: enumeration of quandles

**Theorem:** [Alexander Stein, 2001]

If  $Q$  is a finite latin quandle, then  $\text{LMlt}(Q)$  is solvable.

**Corollary:** Finite latin quandles are solvable.

## Application: enumeration of quandles

**Theorem:** [Alexander Stein, 2001]

If  $Q$  is a finite latin quandle, then  $\text{LMlt}(Q)$  is solvable.

**Corollary:** Finite latin quandles are solvable.

**Theorem:** [Sherman Stein, 1957]

There are **no** latin quandles of order  $\equiv 2 \pmod{4}$ .

**Proof:**

1) Assume it is simple. Then

solvable  $\Rightarrow$  abelian  $\Rightarrow$  affine  $\Rightarrow$  order  $p^k \Rightarrow$  contradiction [Joyce, 1982]

2) Take the smallest counterexample, take a non-trivial congruence, either the factor or a block are smaller of this order.

## Application: enumeration of quandles

**Theorem:** [Alexander Stein, 2001]

If  $Q$  is a finite latin quandle, then  $\text{LMlt}(Q)$  is solvable.

**Corollary:** Finite latin quandles are solvable.

**Theorem:** [Sherman Stein, 1957]

There are **no** latin quandles of order  $\equiv 2 \pmod{4}$ .

**Proof:**

1) Assume it is simple. Then

solvable  $\Rightarrow$  abelian  $\Rightarrow$  affine  $\Rightarrow$  order  $p^k \Rightarrow$  contradiction [Joyce, 1982]

2) Take the smallest counterexample, take a non-trivial congruence, either the factor or a block are smaller of this order.

**More complicated results:** [Bianco, Bonatto, around 2020]

Classification of latin quandles of order  $pq$ , connected quandles of order  $p^3, \dots$

## Application: coloring knots by latin quandles

**Theorem:** [Alexander Stein, 2001]

If  $Q$  is a finite latin quandle, then  $\text{LMlt}(Q)$  is solvable.

**Corollary:** Finite latin quandles are solvable.

**Corollary:** Knots with trivial Alexander polynomial are not colorable by any finite latin quandle.

**Proof idea:**

(1) Bae's theorem: a knot admits a non-trivial coloring by an affine quandle if and only if its Alexander polynomial is non-trivial.

(2) Lemma: If  $c$  is a non-trivial coloring of  $K$  by a quandle  $Q$ , and  $Q = \langle \text{Im}(c) \rangle$ , then  $K$  is colorable by every simple factor of  $Q$ .

$\rightsquigarrow$  If  $Q$  is finite latin, it is solvable, hence all simple factors are affine, and the two facts contradict.

# LOOPS



## Loops

= “non-associative groups”

=  $(Q, \cdot, /, \backslash, 1)$  such that 1 is a unit,  $/, \backslash$  are division operations wrt  $\cdot$

*Translations:*

$$L_x(y) = x \cdot y, \quad R_x(y) = y \cdot x$$

*Multiplication group:*

$$\text{Mlt}(Q) = \langle L_x, R_x : x \in Q \rangle \leq S_Q$$

*Inner mapping group:*

$$\text{Inn}(Q) = (\text{Mlt}(Q))_1 = \langle L_{x,y}, R_{x,y}, T_x : x, y \in Q \rangle \leq \text{Mlt}(Q)$$

where

$$L_{x,y}(z) = (xy) \backslash x(yz), \quad R_{x,y}(z) = (zy)x / (yx), \quad T_x(z) = x \backslash zx$$

## “Naive” commutator theory

[Albert, Bruck, 1940s]

Easy fact: *congruences* correspond to *normal subloops*, i.e., subloops invariant with respect to action of  $\text{Inn}(Q)$ .

*The center:*

$$Z(Q) = \text{Fix}(\text{Inn}(Q)) = \\ = \{a : ax = xa, a(xy) = (ax)y, x(ay) = (xa)y, x(ya) = (xy)a \text{ for all } x, y\}.$$

A loop  $Q$  is called (*classically*) *solvable*, resp. *nilpotent*, if there are normal subloops  $N_i$  such that

$$1 = N_0 \leq N_1 \leq \dots \leq N_k = Q$$

and  $N_{i+1}/N_i$  is an *abelian group*, resp. *contained in the center*  $Z(Q/N_i)$ , for all  $i$ .

# “The” commutator theory

[S, Vojtěchovský, 2014] [Barnes, 2023]

Main theorem:

$$[A, B] = \text{Ng}(I_{u_1, u_2}(a) \setminus I_{v_1, v_2}(a) : I \in \{L, R, T\}, u_i \setminus v_i \in B, a \in A).$$

Consequently,

- $[N, Q] = 1$  iff  $N \leq Z(Q)$ ,
- $[N, N] = 1$  iff  $I_{u_1, u_2}|_N = I_{v_1, v_2}|_N$  for every  $I \in \{L, R, T\}, u_i \setminus v_i \in B$ .
- $[N, N] = 1$  iff  $\varphi|_N \in \text{Aut}(N)$  for every  $\varphi \in \text{Inn}(Q)$  and certain commutators/associators vanish.

Hence, comparing to the “naive” definitions,

- centrality and nilpotence agree,
- abelianness and solvability disagree.

[Drápal 2023] solvability agrees in Moufang loops

# Nilpotent loops

- $|Q| = p^k \Rightarrow Q$  is nilpotent
  - true for groups [quite easy]
  - true for Moufang loops [quite difficult, Glaubermann, Wright 1968]
  - false in general: loops of order  $p$  are abelian groups or counterexamples
- $Q$  finite nilpotent  $\Rightarrow Q \simeq \prod Q_p$  where  $Q_p$  are nilp. loops of order  $p^k$ 
  - true for groups [not difficult]
  - true for Moufang loops [quite difficult, Glaubermann, Wright 1968]
  - false in general: a loop of order 6 is nilpotent, directly indecomposable

$\Rightarrow$  supernilpotence ???

## Higher commutator, supernilpotence

$C_n(\alpha_1, \dots, \alpha_{n-1}; \beta; \gamma)$  iff for every term  $t$  and every  $\bar{a}_i \stackrel{\alpha_i}{\equiv} \bar{b}_i$ ,  $\bar{u} \stackrel{\beta}{\equiv} \bar{v}$

$$t(\bar{x}_1, \dots, \bar{x}_n, \bar{u}) \stackrel{\delta}{\equiv} t(\bar{x}_1, \dots, \bar{x}_n, \bar{v}) \quad \forall (\bar{x}_1, \dots, \bar{x}_n) \in \{\bar{a}_1, \bar{b}_1\} \times \dots \times \{\bar{a}_n, \bar{b}_n\} \\ \neq \{(\bar{b}_1, \dots, \bar{b}_n)\}$$

$\Downarrow$

$$t(\bar{b}_1, \dots, \bar{b}_n, \bar{u}) \stackrel{\delta}{\equiv} t(\bar{b}_1, \dots, \bar{b}_n, \bar{v}).$$

The  $n$ -ary commutator  $[\alpha_1, \dots, \alpha_n]$  is the smallest  $\delta$  such that  $C_n(\alpha_1, \dots, \alpha_{n-1}; \alpha_n; \delta)$ .

**Fact:**  $[\alpha_1, \dots, \alpha_n] \geq [\alpha_1, [\alpha_2, [\dots, [\alpha_{n-1}, \alpha_n]]]]$  (in Mal'tsev varieties)

## Higher commutator, supernilpotence

$C_n(\alpha_1, \dots, \alpha_{n-1}; \beta; \gamma)$  iff for every term  $t$  and every  $\bar{a}_i \stackrel{\alpha_i}{\equiv} \bar{b}_i$ ,  $\bar{u} \stackrel{\beta}{\equiv} \bar{v}$

$$t(\bar{x}_1, \dots, \bar{x}_n, \bar{u}) \stackrel{\delta}{\equiv} t(\bar{x}_1, \dots, \bar{x}_n, \bar{v}) \quad \forall (\bar{x}_1, \dots, \bar{x}_n) \in \{\bar{a}_1, \bar{b}_1\} \times \dots \times \{\bar{a}_n, \bar{b}_n\} \\ \neq \{(\bar{b}_1, \dots, \bar{b}_n)\}$$

$\Downarrow$

$$t(\bar{b}_1, \dots, \bar{b}_n, \bar{u}) \stackrel{\delta}{\equiv} t(\bar{b}_1, \dots, \bar{b}_n, \bar{v}).$$

The  $n$ -ary commutator  $[\alpha_1, \dots, \alpha_n]$  is the smallest  $\delta$  such that  $C_n(\alpha_1, \dots, \alpha_{n-1}; \alpha_n; \delta)$ .

**Fact:**  $[\alpha_1, \dots, \alpha_n] \geq [\alpha_1, [\alpha_2, [\dots, [\alpha_{n-1}, \alpha_n]]]]$  (in Mal'tsev varieties)

An algebra is  $k$ -supernilpotent if  $[1_A, \dots, 1_A] = 0_A$ .

# Supernilpotence – a better “definition”

**Theorem:** [Aichinger, Mudrinski, 2010]

In Mal'tsev varieties,

- 1 an algebra is *k-supernilpotent* if and only if all absorbing polynomials of arity  $> k$  are constant.
- 2 a finite algebra is *k-supernilpotent* if and only if  $A \simeq \prod A_p$  where  $A_p$  is a nilpotent algebra of order power of  $p$

A polynomial is *absorbing* if  $p(a_1, \dots, a_n) = 1$  whenever at least one  $a_i = 1$ .

**Examples:**  $[x, y]$ ,  $[x, y, z]$ ,  $L_{x,y}(z)/z$ , ...,  $[xy, u]/([x, u][y, u])$ , ...

# Supernilpotent groups

**Theorem:** [Aichinger, Ecker, 2006; S, Vojtěchovský 2023]

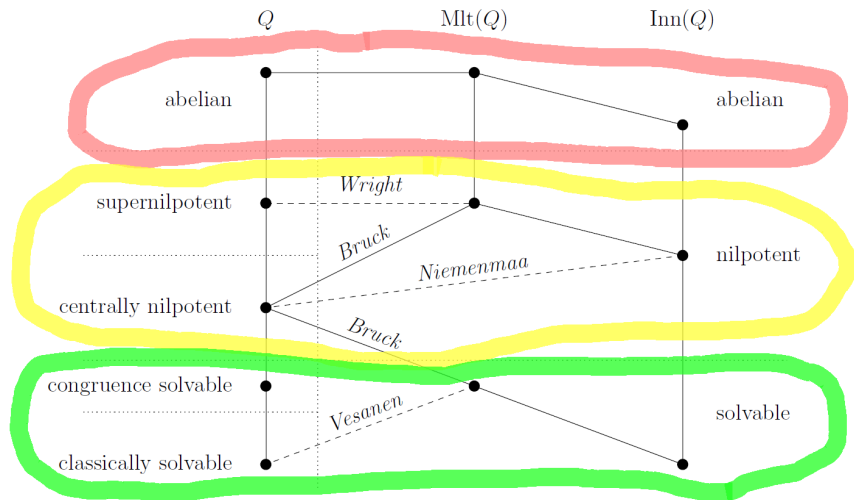
A group is  $k$ -supernilpotent iff  $k$ -nilpotent.

In general, not at all.

- 1  $k$ -supernilpotence  $\Rightarrow$   $k$ -nilpotence
- 2 nilpotence  $\not\Rightarrow$  supernilpotence
- 3 the degree of supernilpotence can be  $\gg$  degree of nilpotence



# Big picture



## Equational basis for 1,2-supernilpotence

Let  $\llbracket x, y \rrbracket$  and  $\llbracket x, y, z \rrbracket$  be any terms such that, in all loops,

$$\llbracket x, y \rrbracket = 1 \Leftrightarrow xy = yx$$

$$\llbracket x, y, z \rrbracket = 1 \Leftrightarrow x(yz) = (xy)z$$

**Example:** the standard commutator and associator

$$\llbracket x, y \rrbracket = (yx) \setminus (xy), \quad \llbracket x, y, z \rrbracket = x(yz) \setminus (xy)z$$

## Equational basis for 1,2-supernilpotence

Let  $\llbracket x, y \rrbracket$  and  $\llbracket x, y, z \rrbracket$  be any terms such that, in all loops,

$$\begin{aligned}\llbracket x, y \rrbracket = 1 &\Leftrightarrow xy = yx \\ \llbracket x, y, z \rrbracket = 1 &\Leftrightarrow x(yz) = (xy)z\end{aligned}$$

**Example:** the standard commutator and associator

$$\llbracket x, y \rrbracket = (yx) \setminus (xy), \quad \llbracket x, y, z \rrbracket = x(yz) \setminus (xy)z$$

**Easy facts:**

1-supernilpotence:  $\llbracket x, y \rrbracket = \llbracket x, y, z \rrbracket = 1$  (abelian groups)

2-supernilpotence:  $\llbracket x, \llbracket y, z \rrbracket \rrbracket = \llbracket x, y, z \rrbracket = 1$  (2-nilpotent groups)

A **group** is  $k$ -nilpotent if and only if  $[x_1, [x_2, [\dots, [x_k, x_{k+1}]]]] = 1$

# Equational basis for 3-supernilpotence

[S, Vojtěchovský, 2023]

TFAE for a loop  $Q$ :

- $Q$  is 3-supernilpotent
- $Q$  satisfies the following identities for all  $[[\cdot, \cdot], [\cdot, \cdot, \cdot]]$
- $Q$  satisfies the following identities for the standard  $[[\cdot, \cdot], [\cdot, \cdot, \cdot]]$

$$1 = [x, [y, u, v]] \tag{1}$$

$$1 = [x, y, [u, v, w]] = [x, [u, v, w], y] = [[u, v, w], x, y] \tag{2}$$

$$1 = [x, y, [u, v]] = [x, [u, v], y] = [[[u, v], x, y]] \tag{3}$$

$$1 = [x, [y, [u, v]]] = [x, [[[u, v], y]]] \tag{4}$$

$$1 = [[[y, [u, v]], x]] = [[[[u, v], y], x]] \tag{5}$$

$$1 = [[[x, y], [u, v]]] \tag{6}$$

$$[xy, u, v] = [x, u, v] [y, u, v] \tag{7}$$

$$[u, xy, v] = [u, x, v] [u, y, v] \tag{8}$$

$$[u, v, xy] = [u, v, x] [u, v, y] \tag{9}$$