Fuzzy logic and the similarity-based interpretation of fuzzy sets

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Our topic:

Reasoning under vagueness

We intend to formalise reasoning about properties which we express in natural language, without using any measuring devices.

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The challenge of vagueness:

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► Combining a coarse and a fine level of granularity: Finding a reasoning system for reasoning with qualitative information in a quantitative framework.

A simple approach

Our quantitative framework: a set of worlds W.



A simple approach

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Our quantitative framework: a set of worlds W.

The qualitative information: for each vague property, its set of **prototypes** in W.

• We consider implications of the form:

 $\alpha_1,\ldots,\alpha_n\to\beta$

"A situation well described by α_1 and ... and α_n is also well described by β ."

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- ▶ These statements cannot be further combined with others.
- $\alpha_1, \ldots, \alpha_n, \beta$ may contain the connectives \land or \lor .
- Contradiction is expressible by the constant \perp .

PGL: Positive Gentzenian Logic

Formulas:

Propositions are built up from variables and \bot, \top by means of \land, \lor . Implications are of the form $\alpha_1, \ldots, \alpha_n \to \beta$.

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Interpretation:

A model for PGL is a non-empty set W, called a set of worlds. An evaluation v maps each proposition to a subset of W, preserving \land, \lor, \bot, \top .

An implication $\alpha \wedge \beta \rightarrow \gamma \vee \delta$ is satisfied by v if

 $v(\alpha) \cap v(\beta) \subseteq v(\gamma) \cup v(\delta).$

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PGL is the Logic of Distributive 0, 1-Lattices.

PGL: a proof system (Font, Verdú)



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A graded approach

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Our quantitative framework: a similarity space (W, s).

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The qualitative information: sets of **prototypes** of vague properties.

Approximate reasoning (RUSPINI)

 $\alpha \xrightarrow{d} \beta$

We consider graded implications:



Connectives in approximate reasoning

The statement

$$\alpha \land \beta \xrightarrow{d} \gamma \lor \delta,$$

reads as:

"If α and β fit to some degree $\geq t \in [0, 1]$, then γ or δ fit to the degree $\geq t \odot d$."

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But how should we interpret the "and" and the "or" here?

Logics for approximate reasoning (Godo, Esteva, Rodríguez, Dubois, Prade, ...)

In a (version of) approximate reasoning, we interpret

$$\alpha \land \beta \xrightarrow{d} \gamma \lor \delta$$

w.r.t. to an evaluation v as

"If a world is similar to $v(\alpha) \cap v(\beta)$ to the degree t, then to $v(\gamma) \cup v(\delta)$ to the degree $\geq t \odot d$."

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Approximate reasoning for reasoning under vagueness

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Benefits:

▶ Clear, transparent framework.

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▶ Clear, transparent framework.

Limitations:

► (Complete) axiomatisation in important cases not known.

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▶ Rules too weak for certain applications.

We want to understand $\alpha \wedge \beta \xrightarrow{d} \gamma \vee \delta$ as

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Then:

- ▶ implicit connection of truth degrees by min/max.
- ▶ much like FL, where degrees are similarities.

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$$U_t(v(\alpha)) \cap U_t(v(\beta)) \subseteq U_{d \odot t}(v(\gamma)) \cup U_{d \odot t}(v(\delta))$$

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for any $t \in [0, 1]$.

Rules for gPGL

$$\begin{split} & \perp \stackrel{d}{\rightarrow} \alpha \qquad \alpha \stackrel{d}{\rightarrow} \alpha \qquad \alpha \stackrel{d}{\rightarrow} \top \qquad \alpha \stackrel{0}{\rightarrow} \beta \\ & \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma, \beta \stackrel{d}{\rightarrow} \alpha} \qquad \frac{\Gamma, \alpha, \beta \stackrel{d}{\rightarrow} \gamma}{\Gamma, \alpha \land \beta \stackrel{d}{\rightarrow} \gamma} \qquad \frac{\Gamma \stackrel{d}{\rightarrow} \alpha \qquad \Gamma \stackrel{d}{\rightarrow} \beta}{\Gamma \stackrel{d}{\rightarrow} \alpha \land \beta} \\ & \frac{\Gamma, \alpha \stackrel{d}{\rightarrow} \gamma \quad \Gamma, \beta \stackrel{d}{\rightarrow} \gamma}{\Gamma, \alpha \lor \beta \stackrel{d}{\rightarrow} \gamma} \qquad \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma \stackrel{d}{\rightarrow} \alpha \lor \beta} \qquad \frac{\Gamma \stackrel{d}{\rightarrow} \beta}{\Gamma \stackrel{d}{\rightarrow} \alpha \lor \beta} \\ & \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma \stackrel{c}{\rightarrow} \alpha} \text{ where } c < d \qquad \frac{\Gamma_{1} \stackrel{c}{\rightarrow} \alpha \quad \Gamma_{2}, \alpha \stackrel{d}{\rightarrow} \beta}{\Gamma_{1}, \Gamma_{2} \stackrel{c \odot d}{\rightarrow} \beta} \end{split}$$

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$$\begin{array}{cccc} \bot \stackrel{d}{\rightarrow} \alpha & \alpha \stackrel{d}{\rightarrow} \alpha & \alpha \stackrel{d}{\rightarrow} \top & \alpha \stackrel{0}{\rightarrow} \beta \\ \\ \hline \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma, \beta \stackrel{d}{\rightarrow} \alpha} & \frac{\Gamma, \alpha, \beta \stackrel{d}{\rightarrow} \gamma}{\Gamma, \alpha \wedge \beta \stackrel{d}{\rightarrow} \gamma} & \frac{\Gamma \stackrel{d}{\rightarrow} \alpha & \Gamma \stackrel{d}{\rightarrow} \beta}{\Gamma \stackrel{d}{\rightarrow} \alpha \wedge \beta} \\ \\ \hline \frac{\Gamma, \alpha \stackrel{d}{\rightarrow} \gamma & \Gamma, \beta \stackrel{d}{\rightarrow} \gamma}{\Gamma, \alpha \vee \beta \stackrel{d}{\rightarrow} \gamma} & \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma \stackrel{d}{\rightarrow} \alpha \vee \beta} & \frac{\Gamma \stackrel{d}{\rightarrow} \beta}{\Gamma \stackrel{d}{\rightarrow} \alpha \vee \beta} \\ \\ \hline \frac{\Gamma \stackrel{d}{\rightarrow} \alpha}{\Gamma \stackrel{c}{\rightarrow} \alpha} \text{ where } c < d & \frac{\Gamma_{1} \stackrel{c}{\rightarrow} \alpha & \Gamma_{2}, \alpha \stackrel{d}{\rightarrow} \beta}{\Gamma_{1}, \Gamma_{2} \stackrel{c \odot d}{\rightarrow} \beta} \end{array}$$

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The rules are well sufficient for practical applications like expert systems.

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Completeness does not hold, however.

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The Graded Positive Gentzenian Logic is "between" FL and AR:

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Furthermore:

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Furthermore:

- ► gPGL is well-applicable, even without a complete proof system.
- ► Completeness *can* be achieved if restricting to "non-disjunctive" theories.

Alternative style of rules for gPGL

gPGL can be axiomatised in alternative ways.



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 gPGL can be axiomatised in alternative ways.

The easiest possibility is to replace

$$\alpha_1,\ldots,\alpha_n \stackrel{d}{\to} \beta$$

by

$$(\alpha_1, s_1), \ldots, (\alpha_n, s_n) \to (\beta, t).$$

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Negation

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Let φ denote vague property. We might want to include $\neg \varphi$ into our language.

By $\neg \varphi$, we mean the negation of φ w.r.t. the coarse level of granularity to which φ refers.

We treat $\neg \varphi$ like φ : $\neg \varphi$ is modelled by the "prototypes" of $\neg \varphi$, that is, by the set of constrasting cases of φ .

Models of negated properties

To interpret \neg , we guess that antitonicity is the minimal assumption.



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Otherwise, we have (at least) the following possibilities of choosing an interpretation $v(\neg \varphi) \subseteq W$:

- We construct $v(\neg \varphi)$ from $v(\varphi)$.
- We construct $v(\neg \varphi)$ from all $v(\psi)$, where ψ contradicts φ .

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- We assume only that $v(\neg \varphi)$ is disjoint from $v(\varphi)$.
- We choose $v(\neg \varphi)$ freely.

Logics with negation

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With the first option, we define nPGL and ngPGL.

For nPGL, v(¬φ) is simply the set-theoretical complement of v(φ).
 This leads to CPL.

Logics with negation

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- For nPGL, v(¬φ) is simply the set-theoretical complement of v(φ). This leads to CPL.
- For ngPGL, the ¹/₂-neighborhoods of v(φ) and v(¬φ) form a partition.
 This corresponds to the standard negation in FL.

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Summary

The Positive Gentzenian Logics ...

- ▶ can be based on rules with a clear intuitive meaning.
- ► are neither fuzzy logics nor logics for approximate reasoning, but something "in between".
- ► are well-applicable, even though completeness can be shown only in a restricted sense.