# Fuzzy Orders for Solving MOLP Problems 

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January 30 - February 03, 2012
FSTA2012

## Problem formulation

In our work we observe a multi-objective linear programming problem, which can be represented as follows:
$\operatorname{MAX} Z$, where $Z=\left(z_{1}, \ldots, z_{k}\right)$ is a vector of objectives,
$z_{i}=\sum_{j=1}^{n} c_{i j} x_{j}$ where $i=1, . ., k$,
subject to $\sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, i=1, \ldots, m$.
That is we should find a vector $x^{0}=\left(x_{1}^{0}, \ldots, x_{n}^{0}\right)$ which maximizes $k$ objective functions with $n$ variables, and $m$ constraints.


$$
\begin{aligned}
& M A X X_{z} \\
& z=\sum_{j=1}^{n} c_{j} x_{j} \\
& \text { subject to } \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i} \\
& i=1, \ldots, m
\end{aligned}
$$



$$
\begin{aligned}
& \operatorname{MAX}\left(z_{1}, z_{2}\right) \\
& z_{i}=\sum_{j=1}^{n} c_{i j} x_{j} \text { where } i=1,2 \\
& \text { subject to } \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, \\
& i=1, \ldots, m .
\end{aligned}
$$

## Fuzzy approach

Membership functions:

$$
\mu_{i}(x)= \begin{cases}0, & z_{i}(x)<z_{i}^{\min } \\ \frac{z_{i}(x)-z_{i}^{\min }}{z_{i}^{\max }-z_{i}^{\min }}, & z_{i}^{\min } \leq z_{i}(x) \leq z_{i}^{\max } \\ 1, & z_{i}(x)>z_{i}^{\max }\end{cases}
$$

$\max _{x} \min _{i} \mu_{i}(x)$


$$
\begin{aligned}
& \operatorname{MAX} z \\
& z=\sum_{j=1}^{n} c_{j} x_{j} \\
& \text { subject to } \sum_{j=1}^{n} a_{i j} x_{j} \leq b_{i}, \\
& i=1, \ldots, m . \\
& x \doteq y \Leftrightarrow z(x)=z(y) \\
& x \preceq y \Leftrightarrow z(x) \leq z(y) \\
& \max (D, \preceq)
\end{aligned}
$$

$$
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& M A X_{n} z \\
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## Fuzzy order approach

1. We define fuzzy order relations $P_{i}$ which generalize the following crisp order relations
$x \preceq_{i} y \Leftrightarrow z_{i}(x) \leq z_{i}(y), i=1, \ldots, k$. Thus each fuzzy order relation describes corresponding objective function $z_{i}$.
2. We aggregate fuzzy orders using an aggregation functior A which preserves the properties of initial fuzzy orders.

$$
P(x, y)=A\left(P_{1}(x, y), \ldots, P_{k}(x, y)\right)
$$

Thus the aggregated fuzzy order relation $P$ provides the information about all objective functions.
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## Definition of a fuzzy equivalence relation

## Definition

A fuzzy binary relation $E: X \times X \rightarrow[0,1]$ on a set $X$ is called fuzzy equivalence relation with respect to a t-norm $T$, for brevity $T$-equivalence, if and only if the following three axioms are fulfilled for all $x, y, z \in X$ :

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1. $E(x, x)=1$ reflexivity;
2. $E(x, y)=E(y, x)$ symmetry;
3. $T(E(x, y), E(y, z)) \leq E(x, z) T$-transitivity.

## Fuzzy equivalence relation

## Theorem

Let $T$ be a continuous Archimedean $t$-norm with an additive generator $t$. For any pseudo-metric $d$, the mapping

$$
E_{d}(x, y)=t^{(-1)}(\min (d(x, y), t(0)))
$$

is a $T$-equivalence.

## Example

Let us consider the set of real numbers $X=\mathbb{R}$ and metric $d(x, y)=|x-y|$ on it. Taking into account that $t_{L}(x)=1-x$ is an additive generator of $T_{L}$ (Łukasiewicz $t$-norm) and that $t_{P}(x)=-\ln (x)$ is an additive generator of $T_{P}$ (product t-norm), we obtain two fuzzy equivalence relations:

$$
\begin{gathered}
E_{L}(x, y)=\max (1-|x-y|, 0) ; \\
E_{P}(x, y)=e^{-|x-y|} .
\end{gathered}
$$

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A pair $(X, P)$ is called a fuzzy ordered set.
A fuzzy ordering $P$ is called strongly linear if and only if

$$
\forall x, y \in X: \quad \max (P(x, y), P(y, x))=1
$$

## Construction of fuzzy orderings (U.Bodenhofer)

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$$
P(x, y)= \begin{cases}1, & \text { if } x \leq y \\ E(x, y), & \text { otherwise }\end{cases}
$$

a strongly linear $T$ - $E$-ordering on $X$.

## Construction of fuzzy orderings



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1. $d_{i}(x, y)=\frac{\left|z_{i}(x)-z_{i}(y)\right|}{z_{i}^{\text {max }}-z_{i}^{\text {min }}}$ are pseudo-metrics
2. $E_{i}(x, y)=t^{(-1)}\left(\min \left(\frac{\left|z_{i}(x)-z_{i}(y)\right|}{z^{\text {max }}-z^{\text {min }}}, t(0)\right)\right)$ are fuzzy $T$-equivalence relations
3. we build fuzzy order relations (T-E-orders) by the following way:
4. $d_{i}(x, y)=\frac{\left|z_{i}(x)-z_{i}(y)\right|}{z_{i}^{\text {max }}-z_{i}^{\text {min }}}$ are pseudo-metrics
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## Example

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3. fuzzy order relations ( $T_{P}$ - $E_{i}$-orders):

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## Aggregation of fuzzy relations

We aggregate fuzzy orders $P_{i} i \in\{1, \ldots, k\}$ using an aggregation function $A$ which preserves the properties of initial fuzzy orders:

$$
P(x, y)=A\left(P_{1}(x, y), \ldots, P_{k}(x, y)\right)
$$

## Definition

Consider an $n$-argument aggregation function $A^{n}:[0,1]^{n} \rightarrow[0,1]$ and an $m$-argument aggregation function $B^{m}:[0,1]^{m} \rightarrow[0,1]$. We say that $A^{n}$ dominates $B^{m}$ if for all $x_{i, j} \in[0,1]$ with $i \in\{1, \ldots, m\}$ and $j \in\{1, \ldots, n\}$ the following property holds:

$$
\begin{aligned}
& B^{m}\left(A^{n}\left(x_{1,1}, \ldots, x_{1, n}\right), \ldots, A^{n}\left(x_{m, 1}, \ldots, x_{m, n}\right)\right) \leq \\
& \leq A^{n}\left(B^{m}\left(x_{1,1}, \ldots, x_{m, 1}\right), \ldots, B^{m}\left(x_{1, n}, \ldots, x_{m, n}\right)\right)
\end{aligned}
$$

## Theorem

Let $|X|>3$ and let $T$ be a t-norm. An aggregation function $A$ preserves $T$-transitivity of fuzzy relations on $X$ if and only if $A$ belongs to the class of aggregation functions which dominate $T$.

## Example

For any $k>2$ and any $p=\left(p_{1}, \ldots, p_{k}\right)$ with $\sum_{i=1}^{k} p_{i} \geq 1$ and
$p_{i} \in[0, \infty] k$-ary aggregation function

$$
A_{p}\left(x_{1}, \ldots, x_{k}\right)=\prod_{i=1}^{k} x_{i}^{p_{i}}
$$

dominates the product t-norm $T_{P}$.

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$$
A_{p}\left(x_{1}, \ldots, x_{k}\right)=\max \left(\sum_{i=1}^{k} x_{i} \cdot p_{i}+1-\sum_{i=1}^{k} p_{i}, 0\right)
$$

dominates the Łukasiewicz t-norm $T_{L}$.

## Theorem

Let $|X|>3$ and let $T$ be a $t$-norm. If $E_{i}$ for all $i \in\{1, \ldots, n\}$ are fuzzy equivalence relations ( $T$-equivalences) then

$$
E(x, y)=A\left(E_{1}(x, y), \ldots, E_{n}(x, y)\right)
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is also a $T$-equivalence relation if $A$ belongs to the class of aggregation functions which dominate $T$.

## Theorem

Let $|X|>3$ and let $T$ be a $t$-norm. If $E_{i}$ for all $i \in\{1, \ldots, n\}$ are fuzzy equivalence relations ( $T$-equivalences); $P_{i}$ for all $i \in\{1, \ldots, n\}$ are fuzzy order relations ( $T$ - $E_{i}$-orders) then
$P(x, y)=A\left(P_{1}(x, y), \ldots, P_{n}(x, y)\right)$ is
$T$ - $E$-order relation if $A$ belongs to the class of aggregation functions which dominate $T$ and

$$
E(x, y)=A\left(E_{1}(x, y), \ldots, E_{n}(x, y)\right) .
$$

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\max _{y} \min _{x} P(x, y)
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## Definition

$x^{*}$ is called Pareto optimal solution if and only if there does not exist another $x \in D$ such that $z_{i}(x) \leq z_{i}\left(x^{*}\right)$ for all $i$ and $z_{j}(x) \neq z_{j}\left(x^{*}\right)$ for at least one $j$.

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## Theorem

An optimal solution $x^{*}$ to the problem $(P)$ is a Pareto optimal solution if it is unique optimal solution.

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An optimal solution $x^{*}$ to the problem $(P)$ is a Pareto optimal solution if for all $x$ and $y z_{i}(x)>z_{i}(y) \Rightarrow P_{i}(x, y)<1, A$ is a strictly monotone function and set $D$ is linearly connected.

## Numerical example

Let us observe the following linear programming problem: $\max z_{1}=x_{1}$,
$\max z_{2}=x_{2}$,
s.t. $x_{1}+x_{2} \leq 1$, $x_{1}, x_{2} \geq 0$.


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## Numerical example(Łukasiewicz t-norm)

We solve the following problem: $\max _{y \in D} \min _{x \in D} P(x, y)$.

$$
f(y)=\min _{x \in B} P(x, y):
$$



Figure:


Figure:

$$
P(x, y)=\frac{P_{1}(x, y)+P_{2}(x, y)}{2}
$$

$$
P(x, y)=\frac{2 P_{1}(x, y)+P_{2}(x, y)}{3}
$$

## Numerical example(Product t-norm)

We solve the following problem: $\max _{y \in D} \min _{x \in D} P(x, y)$.

$$
f(y)=\min _{x \in B} P(x, y):
$$



Figure:


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$$
P(x, y)=P_{1}(x, y)^{\frac{1}{2}} \cdot P_{2}(x, y)^{\frac{1}{2}} \quad P(x, y)=P_{1}(x, y)^{\frac{2}{3}} \cdot P_{2}(x, y)^{\frac{1}{3}}
$$

Thank you for attention!

