

State-morphism MV-algebras and Their Generalizations

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- $\mathcal{S}(M)$ -set of states. $\mathcal{S}(M) \neq \emptyset$.
- **extremal state** $s = \lambda s_1 + (1 - \lambda)s_2$ for $\lambda \in (0, 1) \Rightarrow s = s_1 = s_2$.

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- $s \leftrightarrow \text{Ker}(s)$, 1-1 correspondence

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- μ_s - unique Borel regular σ -additive probability measure on $\mathcal{B}(\mathcal{S}(M))$ such that

$$\mu_s(\partial_e \mathcal{S}(M)) = 1$$

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- MV-algebras with a state are not universal algebras, and therefore, they do not provide an algebraizable logic for probability reasoning over many-valued events
- Flaminio-Montagna - introduce an algebraizable logic whose equivalent algebraic semantics is the variety of state MV-algebras
- A state MV-algebra is a pair (M, τ) , M - MV-algebra, τ unary operation on A s.t.

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- τ -internal operator, state operator

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- state-morphism (M, τ) , τ is an idempotent endomorphism
- s state on M , $[0, 1] \otimes M$,
 $\tau_s(\alpha \otimes a) := \alpha \cdot s(a) \otimes 1$

- $([0, 1] \otimes, \tau_s)$ is an SMV-algebra.

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- if $\tau(M) \in V(S_1, \dots, S_n)$ for some $n \geq 1$, then (M, τ) is an SMMV-algebra
- Iff $\tau((n + 1)x) = \tau(nx)$

Other examples

- Let (G, u) be a unital Riesz space and let $A = \Gamma(G, u)$. Choose an endomorphism $s : A \rightarrow A$ such that $s \circ s = s$ and fix a real number $\lambda \in [0, 1]$. Define a mapping $s_\lambda : A^2 \rightarrow A$ by

$$s_\lambda(x, y) = \lambda s(x) + (1 - \lambda)s(y), \quad (x, y) \in A^2,$$

and set $\tau_{\lambda, s} : A^3 \rightarrow A^3$ by

$$\tau_{\lambda, s}(x, y, z) = (s_\lambda(x, y), s_\lambda(x, y), s_\lambda(x, y)), \quad (x, y, z) \in A^3$$

- Then $D(A)_{\lambda,s} := (A^3, \tau_{\lambda,s})$ is a state MV-algebra such that it is a state morphism MV-algebra iff $\lambda \in \{0, 1\}$.

$$\text{Ker}(\tau_{\lambda,s}) = \begin{cases} \text{Ker}(s) \times \text{Ker}(s) \times A & \text{if } \lambda \in (0, 1), \\ A \times \text{Ker}(s) \times A & \text{if } \lambda = 0, \\ \text{Ker}(s) \times A \times A & \text{if } \lambda = 1. \end{cases}$$

Hence, if A is a subdirectly irreducible MV-algebra and s is faithful, then neither $(A, \tau_{0,s})$ nor $(A, \tau_{1,s})$ is a subdirectly irreducible state morphism MV-algebra.

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- Let $A = [0, 1]$ be the standard MV-algebra and let s be the identity on A . If $\lambda \in \{0, 1\}$, then $D([0, 1]_{1,s})$ generates the variety of state morphism MV-algebras
- Does $D([0, 1]_{\lambda,s})$ for a fixed $\lambda \in (0, 1)$ and the identity s on $[0, 1]$ generate the variety of state MV-algebras ?

State BL-algebras

- M - BL-algebra. A map $\tau : M \rightarrow M$ s.t.

$$(1)_{BL} \tau(0) = 0;$$

$$(2)_{BL} \tau(x \rightarrow y) = \tau(x) \rightarrow \tau(x \wedge y);$$

$$(3)_{BL} \tau(x \odot y) = \tau(x) \odot \tau(x \rightarrow (x \odot y));$$

$$(4)_{BL} \tau(\tau(x) \odot \tau(y)) = \tau(x) \odot \tau(y);$$

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- If $\tau : M \rightarrow M$ is a BL-endomorphism s.t. $\tau \circ \tau = \tau$, - *state-morphism operator* and the couple (M, τ) - *state-morphism BL-algebra*.

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- **Example 0.2** Let M be a BL-algebra. On $M \times M$ we define two operators, τ_1 and τ_2 , as follows

$$\tau_1(a, b) = (a, a), \quad \tau_2(a, b) = (b, b), \quad (a, b) \in M \times M. \quad (2.0)$$

Then τ_1 and τ_2 are two state-morphism operators on $M \times M$.

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- **Example 0.3** Let M be a BL-algebra. On $M \times M$ we define two operators, τ_1 and τ_2 , as follows

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Then τ_1 and τ_2 are two state-morphism operators on $M \times M$.

- $\text{Ker}(\tau) = \{a \in M : \tau(a) = 1\}$.

- We say that two subhoops, A and B , of a BL-algebra M have the *disjunction property* if for all $x \in A$ and $y \in B$, if $x \vee y = 1$, then either $x = 1$ or $y = 1$.

- We say that two subhoops, A and B , of a BL-algebra M have the *disjunction property* if for all $x \in A$ and $y \in B$, if $x \vee y = 1$, then either $x = 1$ or $y = 1$.
- **Lemma 0.5** *Suppose that (M, τ) is a state BL-algebra. Then:*
 - (1) *If τ is faithful, then (M, τ) is a subdirectly irreducible state BL-algebra if and only if $\tau(M)$ is a subdirectly irreducible BL-algebra.*

Now let (M, τ) be subdirectly irreducible.

Then:

- (2) $\text{Ker}(\tau)$ is (either trivial or) a subdirectly irreducible hoop.
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- **Theorem 0.7** *Let (M, τ) be a state BL-algebra satisfying conditions (1), (2) and (3) in the last Lemma. Then (M, τ) is subdirectly irreducible.*

- **Theorem 0.8** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*

- **Theorem 0.9** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*
- (i) M is linear, $\tau = \text{id}_M$, and the BL-reduct M is a subdirectly irreducible BL-algebra.

- **Theorem 0.10** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*
- (i) M is linear, $\tau = \text{id}_M$, and the BL-reduct M is a subdirectly irreducible BL-algebra.
- (ii) The state-morphism operator τ is not faithful, M has no nontrivial Boolean elements, and the BL-reduct M of (M, τ) is a local BL-algebra, $\text{Ker}(\tau)$ is a subdirectly irreducible irreducible hoop, and $\text{Ker}(\tau)$ and $\tau(M)$ have the disjunction property.

- **Theorem 0.11** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*
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- **Theorem 0.12** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*
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- **Theorem 0.13** *A state-morphism BL-algebra (M, τ) is subdirectly irreducible irreducible if and only if one of the following three possibilities holds.*
- (i) M is linear, $\tau = \text{id}_M$, and the BL-reduct M is a subdirectly irreducible BL-algebra.
- (ii) The state-morphism operator τ is not faithful, M has no nontrivial Boolean elements, and the BL-reduct M of (M, τ) is a local BL-algebra, $\text{Ker}(\tau)$ is a subdirectly irreducible irreducible hoop, and $\text{Ker}(\tau)$ and $\tau(M)$ have the disjunction property.

- Moreover, M is linearly ordered if and only if $\text{Rad}_1(M)$ is linearly ordered, and in such a case, M is a subdirectly irreducible BL-algebra such that if F is the smallest nontrivial state-filter for (M, τ) , then F is the smallest nontrivial BL-filter for M .

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- If $\text{Rad}(M) = \text{Ker}(\tau)$, then M is linearly ordered.

- (iii) The state-morphism operator τ is not faithful, M has a nontrivial Boolean element. There are a linearly ordered BL-algebra A , a subdirectly irreducible BL-algebra B , and an injective BL-homomorphism $h : A \rightarrow B$ such that (M, τ) is isomorphic as a state-morphism BL-algebra with the state-morphism BL-algebra $(A \times B, \tau_h)$, where $\tau_h(x, y) = (x, h(x))$ for any $(x, y) \in A \times B$.

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- $D(M) := (M \times M, \tau_M)$
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- $SMMV = V(D([0, 1]))$
- $\mathcal{P}_\tau = V(D(C))$, \mathcal{P} perfect MV-algebras, C-Chang

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$$(\tau(x) \leftrightarrow x)^* \leq (\tau(x) \leftrightarrow x).$$

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- (1) either $n = 0$ or $\text{g.c.d}(n, m) = 1$
- $\forall p \in X, p$ does not divide m
- $\tau(x) =$ standard part of x
- $(A(X), \tau)$ is linearly ordered SMMV-algebra

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- Theorem: Between $\mathcal{MV}\mathcal{I}$ and $\mathcal{MV}\mathcal{R}$ there is uncountably many varieties

Generators of SMBL-algebras

- *t*-norm- function $t : [0, 1] \times [0, 1] \rightarrow [0, 1]$ such that (i) *t* is commutative, associative, (ii) $t(x, 1) = x, x \in [0, 1]$, and (iii) *t* is nondecreasing.

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- If t is continuous, we define $x \odot_t y = t(x, y)$ and $x \rightarrow_t y = \sup\{z \in [0, 1] : t(z, x) \leq y\}$ for $x, y \in [0, 1]$, then $\mathbb{I}_t := ([0, 1], \min, \max, \odot_t, \rightarrow_t, 0, 1)$ is a BL-algebra.

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- The variety of all BL-algebras is generated by all \mathbb{I}_t with a continuous t -norm t .

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- \mathcal{T} denotes the system of all BL-algebras $D(\mathbb{I}_t)$, where t is a continuous t-norm on the interval $[0, 1]$,
- **Theorem 0.15** *The variety of all state-morphism BL-algebras is generated by the class \mathcal{T} .*

General Approach - State-Morphism Algebras

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- **Lemma:** For any $\phi \in \text{Con } \tau(\mathbf{A})$, we have $\theta_\phi \in \text{Con } (\mathbf{A}, \tau)$, and $\theta_\phi \cap \tau(A)^2 = \phi$. In addition, $\theta_\tau \in \text{Con } (\mathbf{A}, \tau)$, $\phi \subseteq \theta_\phi$, and $\Theta_\tau(\phi) \subseteq \theta_\phi$.

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- **Lemma:** Let $\theta \in \text{Con } \mathbf{A}$ be such that $\theta \subseteq \theta_\tau$. Then $\theta \in \text{Con } (\mathbf{A}, \tau)$ holds.
- **Lemma:** If $x, y \in \tau(\mathbf{A})$, then $\Theta(x, y) = \Theta_\tau(x, y)$. Consequently, $\Theta(\phi) = \Theta_\tau(\phi)$ whenever $\phi \subseteq \tau(A)^2$.
- if $(\mathbf{C}, \tau) \hookrightarrow (\mathbf{B} \times \mathbf{B}, \tau_B)$, (\mathbf{C}, τ) is said to be a *subdiagonal state-morphism algebra*

- **Theorem 0.16** *Let (\mathbf{A}, τ) be a subdirectly irreducible state-morphism algebra such that \mathbf{A} is subdirectly reducible. Then there is a subdirectly irreducible algebra \mathbf{B} such that (\mathbf{A}, τ) is \mathbf{B} -subdiagonal.*

- **Theorem 0.18** *Let (\mathbf{A}, τ) be a subdirectly irreducible state-morphism algebra such that \mathbf{A} is subdirectly reducible. Then there is a subdirectly irreducible algebra \mathbf{B} such that (\mathbf{A}, τ) is \mathbf{B} -subdiagonal.*
- **Theorem 0.19** *For every subdirectly irreducible state-morphism algebra (\mathbf{A}, τ) , there is a subdirectly irreducible algebra \mathbf{B} such that (\mathbf{A}, τ) is \mathbf{B} -subdiagonal.*

- **Theorem 0.20** *Let (\mathbf{A}, τ) be a subdirectly irreducible state-morphism algebra such that \mathbf{A} is subdirectly reducible. Then there is a subdirectly irreducible algebra \mathbf{B} such that (\mathbf{A}, τ) is \mathbf{B} -subdiagonal.*
- **Theorem 0.21** *For every subdirectly irreducible state-morphism algebra (\mathbf{A}, τ) , there is a subdirectly irreducible algebra \mathbf{B} such that (\mathbf{A}, τ) is \mathbf{B} -subdiagonal.*
- \mathcal{K} of algebras of the same type, $I(\mathcal{K})$, $H(\mathcal{K})$, $S(\mathcal{K})$ and $P(\mathcal{K})$ $D(\mathcal{K})$

- **Theorem 0.22** (1) *For every class \mathcal{K} of algebras of the same type F ,
 $V(D(\mathcal{K})) = V(\mathcal{K})_\tau$.*
(2) *Let \mathcal{K}_1 and \mathcal{K}_2 be two classes of same type algebras. Then $V(D(\mathcal{K}_1)) = V(D(\mathcal{K}_2))$ if and only if $V(\mathcal{K}_1) = V(\mathcal{K}_2)$.*

- Theorem 0.24** (1) *For every class \mathcal{K} of algebras of the same type F ,*

$$V(D(\mathcal{K})) = V(\mathcal{K})_\tau.$$
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- Theorem 0.25** *If a system \mathcal{K} of algebras of the same type F generates the whole variety $\mathcal{V}(F)$ of all algebras of type F , then the variety $\mathcal{V}(F)_\tau$ of all state-morphism algebras (\mathbf{A}, τ) , where $\mathbf{A} \in \mathcal{V}(F)$, is generated by the class $\{D(\mathbf{A}) : \mathbf{A} \in \mathcal{K}\}$.*

- **Theorem 0.26** *If \underline{A} is a subdirectly irreducible algebra, then any state-morphism algebra (\underline{A}, τ) is subdirectly irreducible.*

- **Theorem 0.28** *If \underline{A} is a subdirectly irreducible algebra, then any state-morphism algebra (\underline{A}, τ) is subdirectly irreducible.*
- **Theorem 0.29** *A variety \mathcal{V}_τ satisfy the CEP if and only if \mathcal{V} satisfies the CEP.*

Applications

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- The variety of all state-morphism MTL-algebras is generated by the class $\{D(\mathbb{I}_t) : \mathbb{I}_t \in \mathcal{T}_{lc}\}$.

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- The variety of all state-morphism naBL-algebras is generated by the class $\{D(\mathbb{I}_t^{na}) : \mathbb{I}_t \in na\mathcal{T}\}$.

- If a unital ℓ -group (G, u) is double transitive, then $D(\Gamma(G, u))$ generates the variety of state-morphism pseudo MV-algebras.

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Thank you for your attention