# On multi-objective linear programming approach for solving fuzzy matrix games 

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Eleventh international conference on fuzzy set theory and applications

Liptovsky Jan, Slovak Republic
January 30 - February 3, 2012

In this talk we deal with non cooperative two-person games with fuzzy pay-offs. Namely, we consider matrix games where each component of the pay-off matrix is a fuzzy number. We describe the formal definition of the value of a fuzzy pay-off matrix game and develop a fuzzy programming method to find it by solving the corresponding bilevel linear programming problem. To realize fuzzy programming on two levels we apply specially designed aggregation of objectives.

## Matrix games

Matrix games are zero - sum two - person games. A is called the pay-off matrix: if Player I chooses $\mathrm{i}^{\text {th }}(\mathrm{i}=\overline{1, \mathrm{~m}})$ strategy and Player II chooses $\mathrm{j}^{\text {th }}(\mathrm{j}=\overline{1, \mathrm{n}})$ strategy then $\mathrm{a}_{\mathrm{ij}}$ is amount paid by Player II to Player I:

$$
A=\left(\begin{array}{llll}
a_{11} & a_{12} & \ldots & a_{1 n} \\
a_{21} & a_{22} & \ldots & a_{2 n} \\
\ldots & \ldots & \ldots & \ldots \\
a_{m 1} & a_{m 2} & \ldots & a_{m n}
\end{array}\right)
$$

Lower value of the game: $v=\max _{\mathrm{i}=\overline{1, \mathrm{~m}}} \min _{\mathrm{j}=\overline{1, \mathrm{n}}} \mathrm{a}_{\mathrm{ij}}$.
Upper value of the game: $w=\min _{\mathrm{j}=1, \mathrm{n}} \max _{\mathrm{i}=\overline{1, \mathrm{~m}}} \mathrm{a}_{\mathrm{ij}}$. If $\mathrm{v}=\mathrm{w}$ then it is called value of the game.

Expected pay-off function

$$
\mathrm{E}_{\mathrm{A}}(\mathrm{x}, \mathrm{y})=\sum_{\mathrm{i}=1}^{\mathrm{m}} \sum_{\mathrm{j}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{ij}} \mathrm{x}_{\mathrm{i}} \mathrm{y}_{\mathrm{j}}
$$

when Player I chooses mixed strategy x and Player II - y:

$$
\begin{aligned}
\mathrm{S}^{\mathrm{m}}=\left\{\mathrm{x}=\left(\mathrm{x}_{1}, \mathrm{x}_{2}, \ldots, \mathrm{x}_{\mathrm{m}}\right):\right. & \left.\sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{x}_{\mathrm{i}}=1, \quad \mathrm{x}_{\mathrm{i}} \in[0,1]\right\}, \\
\mathrm{S}^{\mathrm{n}}=\left\{\mathrm{y}=\left(\mathrm{y}_{1}, \mathrm{y}_{2}, \ldots, \mathrm{y}_{\mathrm{n}}\right):\right. & \left.\sum_{\mathrm{j}=1}^{\mathrm{n}} \mathrm{y}_{\mathrm{j}}=1, \quad \mathrm{y}_{\mathrm{j}} \in[0,1]\right\} .
\end{aligned}
$$

The function $\mathrm{E}_{\mathrm{A}}$ has a seddle point $\left(\mathrm{x}^{*}, \mathrm{y}^{*}\right)$ and $\mathrm{x}^{*}, \mathrm{y}^{*}$ are optimal strategies for Player I and Player II.
Value of the game $\mathrm{E}_{\mathrm{A}}\left(\mathrm{x}^{*}, \mathrm{y}^{*}\right)$.

## Linear programming and matrix game equivalence

$$
\begin{aligned}
& \text { Player I: } \\
& \mathrm{v} \rightarrow \max \\
& \qquad \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{a}_{\mathrm{ij}} \mathrm{x}_{\mathrm{i}} \geqslant \mathrm{v} \quad(\mathrm{j}=\overline{1, \mathrm{n}}) \\
& \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{x}_{\mathrm{i}}=1 \\
& \mathrm{x}_{\mathrm{i}} \geqslant 0 \quad(\mathrm{i}=\overline{1, \mathrm{~m}})
\end{aligned}
$$

## Player II:

$$
\mathrm{w} \rightarrow \min ,
$$

$$
\sum_{\mathrm{j}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{ij}} \mathrm{y}_{\mathrm{j}} \leqslant \mathrm{w} \quad(\mathrm{i}=\overline{1, \mathrm{~m}})
$$

$$
\sum_{j=1}^{n} y_{j}=1
$$

$$
\mathrm{y}_{\mathrm{j}} \geqslant 0 \quad(\mathrm{j}=\overline{1, \mathrm{n}})
$$

Solutions of the problems: $\mathrm{x}^{*}, \mathrm{y}^{*}$ and $\mathrm{v}^{*}, \mathrm{w}^{*}$.

## Fuzzy matrix games

Pay-off matrix:

$$
\tilde{A}=\left(\begin{array}{llll}
\tilde{a}_{11} & \tilde{a}_{12} & \ldots & \tilde{a}_{1 n} \\
\tilde{a}_{21} & \tilde{a}_{22} & \ldots & \tilde{a}_{2 n} \\
\ldots & \ldots & \ldots & \ldots \\
\tilde{a}_{m 1} & \tilde{a}_{m 2} & \ldots & \tilde{a}_{m n}
\end{array}\right)
$$

where $\tilde{\mathrm{a}}_{\mathrm{ij}}(\mathrm{i}=\overline{1, \mathrm{~m}} ; \mathrm{j}=\overline{1, \mathrm{n}})$ are fuzzy numbers.

## Fuzzy numbers (FN)

Fuzzy number is a function $\tilde{a}_{\mathrm{ij}}: \mathrm{R} \rightarrow[0,1]$ if:

- there exists the unique point $\mathrm{a}_{\mathrm{ij}}^{\mathrm{M}}$ such that $\tilde{\mathrm{a}}_{\mathrm{ij}}\left(\mathrm{a}_{\mathrm{ij}}^{\mathrm{M}}\right)=1$;
- $\alpha$ - cuts $\left.\tilde{\mathrm{a}}_{\mathrm{ij}}\right|_{\alpha}$ are closed for all $\alpha \in[0,1]$.


We denote: $\tilde{a}_{\mathrm{ij}}^{\mathrm{L}}=\left.\tilde{a}_{\mathrm{ij}}\right|_{\left.]-\infty ; \mathrm{a}_{\mathrm{ij}}^{\mathrm{M}}\right]}, \tilde{a}_{\mathrm{ij}}^{\mathrm{U}}=\left.\tilde{\mathrm{a}}_{\mathrm{ij}}\right|_{\left[\mathrm{a}_{\mathrm{ij}}^{\mathrm{M}} ; \infty[ \right.}[$.
$\alpha$ - cut for $\tilde{\mathrm{a}}_{\mathrm{ij}}$ will be $\left[\left.\mathrm{a}_{\mathrm{ij}}^{\mathrm{L}}\right|_{\alpha} ;\left.\mathrm{a}_{\mathrm{ij}}^{\mathrm{U}}\right|_{\alpha}\right]$.

## Triangular fuzzy numbers (TFN)

Triangular fuzzy number:

$$
\tilde{a}_{i j}=\left(a_{i j}^{L}, a_{i j}^{M}, a_{i j}^{U}\right), \text { where } a_{i j}^{L}<a_{i j}^{M}<a_{i j}^{U} .
$$



Membership function:

$$
\tilde{a}_{i j}\left(a_{i j}\right)= \begin{cases}0, & \text { if } a_{i j} \leqslant a_{i j}^{L}, a_{i j} \geqslant a_{i j}^{U}, \\ \frac{a_{i j}-a_{i j}^{L}}{a_{i j}^{M}-a_{i j}^{L}}, & \text { if } a_{i j}^{L} \leqslant a_{i j} \leqslant a_{i j}^{M}, \\ \frac{i_{i j}-a_{i j}^{U}}{a_{i j}^{M}-a_{i j}^{U}}, & \text { if } a_{i j}^{M} \leqslant a_{i j} \leqslant a_{i j}^{U} .\end{cases}
$$

## Multi-objective programming approach: Li's model

$$
\begin{aligned}
& \text { Player } 1 \\
& \left(v^{L}, v^{M}, v^{U}\right) \rightarrow \max \\
& \qquad \sum_{i=1}^{m} a_{i j}^{L} x_{i} \geqslant v^{L}, \quad(j=1, . ., n) \\
& \sum_{i=1}^{m} a_{i j}^{M} x_{i} \geqslant v^{M}, \quad(j=1, \ldots, n) \\
& \sum_{i=1}^{m} a_{i j}^{U} x_{i} \geqslant v^{U}, \quad(j=1, \ldots, n) \\
& \sum_{i=1}^{m} x_{i}=1, x_{i} \geqslant 0
\end{aligned}
$$

## Reasonable solution of fuzzy matrix games

## Ordering of TFN's

$\tilde{\mathrm{t}}=\left(\mathrm{t}^{\mathrm{L}}, \mathrm{t}^{\mathrm{M}}, \mathrm{t}^{\mathrm{U}}\right), \tilde{\tau}=\left(\tau^{\mathrm{L}}, \tau^{\mathrm{M}}, \tau^{\mathrm{U}}\right)$ are TFN's. Then $\tilde{\mathrm{t}} \leqslant \tilde{\tau}$ if
$\mathrm{t}^{\mathrm{L}} \leqslant \tau^{\mathrm{L}}, \mathrm{t}^{\mathrm{M}} \leqslant \tau^{\mathrm{M}}$ and $\mathrm{t}^{\mathrm{U}} \leqslant \tau^{\mathrm{U}}$.

## Definition

Let $\tilde{\mathrm{v}}=\left(\mathrm{v}^{\mathrm{L}}, \mathrm{v}^{\mathrm{M}}, \mathrm{v}^{\mathrm{U}}\right)$ and $\tilde{\mathrm{w}}=\left(\mathrm{w}^{\mathrm{L}}, \mathrm{w}^{\mathrm{M}}, \mathrm{w}^{\mathrm{U}}\right)$ be TFN's. Then $(\tilde{\mathrm{v}}, \tilde{\mathrm{w}})$ is called a reasonable solution of the fuzzy matrix game if there exist $\bar{x} \in S^{m}, \bar{y} \in S^{n}$ such that

- $\mathrm{E}_{\tilde{\mathrm{A}}}(\overline{\mathrm{x}}, \mathrm{y}) \geqslant \tilde{\mathrm{v}}$ for all $\mathrm{y} \in \mathrm{S}^{\mathrm{n}}$;
- $\mathrm{E}_{\tilde{\mathrm{A}}}(\mathrm{x}, \overline{\mathrm{y}}) \leqslant \tilde{\mathrm{w}}$ for all $\mathrm{x} \in \mathrm{S}^{m}$.

If ( $\tilde{\mathrm{v}}, \tilde{\mathrm{w}}$ ) is a reasonable solution of fuzzy game then $\tilde{\mathrm{v}}$ (respectively $\tilde{w}$ ) is called the reasonable value of Player I (Player II).

## Solutions of fuzzy matrix games

V ( respectively W ) is the set of all reasonable values ṽ ( respectively w $)$ for Player I (Player II).

## Definition

An element $\left(\tilde{v}_{*}=\left(v_{*}^{\mathrm{L}}, \mathrm{v}_{*}^{\mathrm{M}}, \mathrm{v}_{*}^{\mathrm{U}}\right), \tilde{\mathrm{w}}_{*}=\left(\mathrm{w}_{*}^{\mathrm{L}}, \mathrm{w}_{*}^{\mathrm{M}}, \mathrm{w}_{*}^{\mathrm{U}}\right)\right) \in \mathrm{V} \times \mathrm{W}$ is called a solution of the fuzzy game if

- there does not exist any $\tilde{\mathrm{v}}=\left(\mathrm{v}^{\mathrm{L}}, \mathrm{v}^{\mathrm{M}}, \mathrm{v}^{\mathrm{U}}\right) \in \mathrm{V}$ such that $\left(\mathrm{v}^{\mathrm{L}}, \mathrm{v}^{\mathrm{M}}, \mathrm{v}^{\mathrm{U}}\right) \geqslant\left(\mathrm{v}_{*}^{\mathrm{L}}, \mathrm{v}_{*}^{\mathrm{M}}, \mathrm{v}_{*}^{\mathrm{U}}\right)$;
- there does not exist any $\tilde{\mathrm{w}}=\left(\mathrm{w}^{\mathrm{L}}, \mathrm{w}^{\mathrm{M}}, \mathrm{w}^{\mathrm{U}}\right) \in \mathrm{W}$ such that $\left(\mathrm{w}^{\mathrm{L}}, \mathrm{w}^{\mathrm{M}}, \mathrm{w}^{\mathrm{U}}\right) \leqslant\left(\mathrm{w}_{*}^{\mathrm{L}}, \mathrm{w}_{*}^{\mathrm{M}}, \mathrm{w}_{*}^{\mathrm{U}}\right)$.


## Step 1

$\mathrm{v}^{\mathrm{M}} \rightarrow \max _{\mathrm{x} \in \mathrm{D}}$
Solution: $\mathrm{v}_{*}^{\mathrm{M}}$ and $\mathrm{x}^{*}$
Step 2
$\mathrm{v}_{\mathrm{m}}^{\mathrm{L}}, \mathrm{v}^{\mathrm{U}} \rightarrow \max$
$\sum_{i=1} a_{i j}^{L} x_{i}^{*} \geqslant v^{L}, \quad(j=1, . ., n)$
m
$\sum_{i=1} a_{i j}^{U} x_{i}^{*} \geqslant v^{U}, \quad(j=1, . ., n)$
Solution of the problem is: $\left(\mathrm{v}_{*}^{\mathrm{L}}, \mathrm{v}_{*}^{\mathrm{M}}, \mathrm{v}_{*}^{\mathrm{U}}\right)$.

## Numerical example 1

$$
\tilde{\mathrm{A}}=\left(\begin{array}{cc}
\tilde{20} & \tilde{5} \\
\tilde{10} & \tilde{20}
\end{array}\right),
$$

where $\tilde{20}=(10,20,30), \tilde{5}=(1,5,20), \tilde{10}=(9,10,15)$,
$\tilde{20}=(2,20,30)$ are TFN's.
If we solve individual problems of

$$
\left(\mathrm{v}^{\mathrm{L}}, \mathrm{v}^{\mathrm{M}}, \mathrm{v}^{\mathrm{U}}\right) \rightarrow \max _{\mathrm{x} \in \mathrm{D}_{1}}
$$

then

$$
\mathrm{x}_{*}^{\mathrm{M}}=(0.4,0.6), \mathrm{x}_{*}^{\mathrm{L}}=(0,1), \mathrm{x}_{*}^{\mathrm{U}}=(0.6,0.4)
$$

## Bi-level linear programming

## Player 1

$$
\begin{aligned}
& P_{1}^{1}: v^{M} \rightarrow \max \\
& P_{2}^{1}: v^{L} \rightarrow \max \\
& v^{\mathrm{U}} \rightarrow \max \\
& \quad \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{L}} \mathrm{x}_{\mathrm{i}} \geqslant \mathrm{v}^{\mathrm{L}} \quad(\mathrm{j}=\overline{1, \mathrm{n}}), \\
& \quad \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{M}} \mathrm{x}_{\mathrm{i}} \geqslant \mathrm{v}^{\mathrm{M}} \quad(\mathrm{j}=\overline{1, \mathrm{n}}), \\
& \quad \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{U}} \mathrm{x}_{\mathrm{i}} \geqslant \mathrm{v}^{\mathrm{U}} \quad(\mathrm{j}=\overline{1, \mathrm{n}}), \\
& \quad \sum_{\mathrm{i}=1}^{\mathrm{m}} \mathrm{x}_{\mathrm{i}}=1, \mathrm{x}_{\mathrm{i}} \geqslant 0 .
\end{aligned}
$$

## Player 2

$$
\begin{gathered}
\mathrm{P}_{1}^{2}: \mathrm{w}^{\mathrm{M}} \rightarrow \min \\
\mathrm{P}_{2}^{2}: \mathrm{w}^{\mathrm{L}} \rightarrow \min \\
\mathrm{w}^{\mathrm{U}} \rightarrow \min
\end{gathered}
$$

$$
\sum_{\mathrm{j}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{L}} \mathrm{y}_{\mathrm{j}} \leqslant \mathrm{w}^{\mathrm{L}} \quad(\mathrm{i}=\overline{1, \mathrm{~m}})
$$

$$
\sum_{j=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{M}} \mathrm{y}_{\mathrm{j}} \leqslant \mathrm{w}^{\mathrm{M}} \quad(\mathrm{i}=\overline{1, \mathrm{~m}})
$$

$$
\sum_{\mathrm{j}=1}^{\mathrm{n}} \mathrm{a}_{\mathrm{ij}}^{\mathrm{U}} \mathrm{y}_{\mathrm{j}} \leqslant \mathrm{w}^{\mathrm{U}} \quad(\mathrm{i}=\overline{1, \mathrm{~m}})
$$

$$
\sum_{j=1}^{n} y_{j}=1, y_{j} \geqslant 0
$$

## Individual solutions

$\square \mathrm{v}^{\mathrm{M}} \rightarrow \max$
$\mathrm{D}_{1}$
Solutions:

$$
v_{\max }^{\mathrm{M}}=\mathrm{v}^{\mathrm{M}}\left(\mathrm{x}^{\mathrm{M}}\right), \text { where } \mathrm{x}^{\mathrm{M}}=\left(\mathrm{x}_{1}^{\mathrm{M}}, \ldots, \mathrm{x}_{\mathrm{m}}^{\mathrm{M}}\right)
$$

- $\mathrm{v}^{\mathrm{L}} \rightarrow \max$
$\mathrm{D}_{1}$
Solutions:

$$
\mathrm{v}_{\text {max }}^{\mathrm{L}}=\mathrm{v}^{\mathrm{L}}\left(\mathrm{x}^{\mathrm{L}}\right), \text { where } \mathrm{x}^{\mathrm{L}}=\left(\mathrm{x}_{1}^{\mathrm{L}}, \ldots, \mathrm{x}_{\mathrm{m}}^{\mathrm{L}}\right)
$$

- $\mathrm{v}^{\mathrm{U}} \rightarrow \max _{\mathrm{D}_{1}}$

Solutions:

$$
\mathrm{v}_{\max }^{\mathrm{U}}=\mathrm{v}^{\mathrm{U}}\left(\mathrm{x}^{\mathrm{U}}\right), \text { where } \mathrm{x}^{\mathrm{U}}=\left(\mathrm{x}_{1}^{\mathrm{U}}, \ldots, \mathrm{x}_{\mathrm{m}}^{\mathrm{U}}\right)
$$

## Membership function

$$
\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right)= \begin{cases}0, & \text { if } v^{\mathrm{M}} \leqslant \mathrm{v}_{0}^{\mathrm{M}} \\ \frac{v^{M}-v_{0}^{\mathrm{M}}}{\mathrm{v}_{\max }^{\mathrm{M}}-\mathrm{v}_{0}^{\mathrm{M}}}, & \text { if } \mathrm{v}_{0}^{\mathrm{M}} \leqslant v^{\mathrm{M}} \leqslant \mathrm{v}_{\max }^{\mathrm{M}} \\ 1, & \text { if } v^{\mathrm{M}} \geqslant \mathrm{v}_{\max }^{\mathrm{M}}\end{cases}
$$

where $\mathrm{v}_{0}^{\mathrm{M}}=\min \left(\mathrm{v}^{\mathrm{M}}\left(\mathrm{x}^{\mathrm{L}}\right), \mathrm{v}^{\mathrm{M}}\left(\mathrm{x}^{\mathrm{U}}\right)\right)$.


To maximize membership functions:

$$
\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right), \mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right), \mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right) \longrightarrow \max _{\mathrm{x} \in \mathrm{D}_{1}}
$$

we maximize the smallest extreme degree of achievement among all functions:

$$
\min \left\{\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right), \mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right), \mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right)\right\} \longrightarrow \max _{\mathrm{x} \in \mathrm{D}_{1}}
$$

We denote $\sigma:\left\langle\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right) \geqslant \sigma, \mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right) \geqslant \sigma, \mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right) \geqslant \sigma\right\rangle$ and write problem in this form:

$$
\begin{gathered}
\sigma \longrightarrow \max _{\mathrm{x}, \sigma} \\
\left\{\begin{array}{l}
\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right) \geqslant \sigma \\
\mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right) \geqslant \sigma \\
\mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right) \geqslant \sigma \\
\mathrm{x} \in \mathrm{D}_{1}
\end{array}\right.
\end{gathered}
$$

- membership function $\mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right)$

- membership function $\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right)$

- membership function $\mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right)$



$$
\sigma=\max _{\mathrm{x}: \mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right)=\alpha} \min \left\{\mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right), \mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right)\right\}
$$

This graph is obtained by solving the following problems:

$$
\begin{gathered}
\sigma \longrightarrow \max _{\mathrm{x}, \sigma} \\
\left\{\begin{array}{l}
\mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right)=\alpha \\
\mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right) \geqslant \sigma \\
\mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right) \geqslant \sigma \\
\mathrm{x} \in \mathrm{D}_{1}
\end{array}\right.
\end{gathered}
$$

## Numerical example 2

Numerical example (Li and Yang, Campos)

$$
\tilde{A}=\left(\begin{array}{cc}
1 \tilde{8} 0 & \tilde{90} \\
1 \tilde{5} 6 & 1 \tilde{8} 0
\end{array}\right)
$$

where $1 \tilde{8} 0=(175,180,190), 1 \tilde{5} 6=(150,156,158)$, $\tilde{90}=(80,90,100)$ are TFN's.

Individual solutions:

- $\mathrm{v}^{\mathrm{M}} \rightarrow \max _{\mathrm{D}_{1}}$

Solutions: $\mathrm{v}_{\text {max }}^{\mathrm{M}}=161.052, \mathrm{x}^{\mathrm{M}}=(0.789,0.210)$

- $\mathrm{v}^{\mathrm{L}} \rightarrow \max _{\mathrm{D}_{1}}$

Solutions: $\mathrm{v}_{\text {max }}^{\mathrm{L}}=155.208, \mathrm{x}^{\mathrm{L}}=(0.791,0.208)$

- $\mathrm{v}^{\mathrm{U}} \rightarrow \max _{\mathrm{D}_{1}}$
$\mathrm{D}_{1}$
Solutions: $\mathrm{v}_{\text {max }}^{\mathrm{U}}=166.393, \mathrm{x}^{\mathrm{U}}=(0.737,0.262)$

$$
\mathrm{v}_{0}^{\mathrm{M}}=156.393, \mathrm{v}_{0}^{\mathrm{L}}=150.081, \mathrm{v}_{0}^{\mathrm{U}}=164.666
$$

## Membership functions

$$
\begin{aligned}
& \mu^{\mathrm{M}}\left(\mathrm{v}^{\mathrm{M}}\right)= \begin{cases}0, & \text { if } \mathrm{v}^{\mathrm{M}} \leqslant 156.393 \\
\frac{v^{\mathrm{M}}-156.393}{161.052-156.393}, & \text { if } 156.393 \leqslant \mathrm{v}^{\mathrm{M}} \leqslant 161.052 \\
1, & \text { if } \mathrm{v}^{\mathrm{M}} \geqslant 161.052\end{cases} \\
& \mu^{\mathrm{L}}\left(\mathrm{v}^{\mathrm{L}}\right)= \begin{cases}0, & \text { if } \mathrm{v}^{\mathrm{L}} \leqslant 150.081 \\
\frac{\mathrm{v}^{\mathrm{L}}-150.081}{155.208-150.081}, & \text { if } 150.081 \leqslant \mathrm{v}^{\mathrm{L}} \leqslant 155.208 \\
1, & \text { if } \mathrm{v}^{\mathrm{L}} \geqslant 155.208\end{cases} \\
& \mu^{\mathrm{U}}\left(\mathrm{v}^{\mathrm{U}}\right)= \begin{cases}0, & \text { if } \mathrm{v}^{\mathrm{U}} \leqslant 164.666 \\
\frac{\mathrm{v}^{\mathrm{U}}-164.666}{166.393-164.666}, & \text { if } 164.666 \leqslant \mathrm{v}^{\mathrm{U}} \leqslant 166.393 \\
1, & \text { if } \mathrm{v}^{\mathrm{U}} \geqslant 166.393\end{cases}
\end{aligned}
$$

$$
\begin{gathered}
\sigma \longrightarrow \max _{\mathrm{x}, \sigma} \\
\left\{\begin{array}{l}
\frac{\mathrm{v}^{\mathrm{M}}-156.393}{161.052-156.393} \geqslant \sigma \\
\frac{\mathrm{v}^{\mathrm{L}}-150.081}{155.208-150.081} \geqslant \sigma \\
\frac{\mathrm{v}^{\mathrm{U}}-164.666}{166.393-164.666} \geqslant \sigma \\
\mathrm{x} \in \mathrm{D}_{1}
\end{array}\right.
\end{gathered}
$$

Solution:

$$
\sigma^{*}=0.50000001796808
$$

|  | $\alpha$ | $\sigma$ | $\mathrm{v}^{\mathrm{M}}$ |
| :---: | :---: | :---: | :---: |
|  | 0 | 0.500000017968089905 | 156.3934 |
|  | 0.1 | 0.500000017968088351 | 156.8593 |
| $\left\{\begin{array}{l} \frac{\mathrm{v}^{\mathrm{M}}-156.393}{161.052-156.393}=\alpha \\ \frac{\mathrm{v}^{\mathrm{L}}-150.081}{155.208-150.081} \geqslant \sigma \\ \frac{\mathrm{v}^{\mathrm{U}}-164.666}{166.393-164.666} \geqslant \sigma \\ \mathrm{x} \in \mathrm{D}_{1} \end{array}\right.$ | 0.2 | 0.500000017968085908 | 157.3252 |
|  | 0.3 | 0.500000017968091681 | 157.7911 |
|  | 0.4 | 0.500000017968088573 | 158.2571 |
|  | 0.5 | 0.500000017968092902 | 158.7230 |
|  | 0.6 | 0.424383778061159233 | 159.1889 |
|  | 0.7 | 0.328447736942724334 | 159.6548 |
|  | 0.8 | 0.232511695824289433 | 160.1207 |
|  | 0.9 | 0.136575654705872296 | 160.5867 |
|  | 1 | 0.0406396135874516064 | 161.0526 |



If we choose level $\delta=0.8$ then
$\mathrm{x}_{* *}=(0.78,0.22)$
$\mu^{\mathrm{M}}\left(\mathrm{v}_{* *}^{\mathrm{M}}\right)=0.8, \mu^{\mathrm{L}}\left(\mathrm{v}_{* *}^{\mathrm{L}}\right)=0.232511701, \mu^{\mathrm{U}}\left(\mathrm{v}_{* *}^{\mathrm{U}}\right)=0.232511687$.

## Numerical example 1

$$
\tilde{\mathrm{A}}=\left(\begin{array}{cc}
\tilde{20} & \tilde{5} \\
\tilde{10} & \tilde{20}
\end{array}\right),
$$

where $\tilde{20}=(10,20,30), \tilde{5}=(1,5,20), \tilde{10}=(9,10,15)$, $\tilde{20}=(2,20,30)$ are TFN's.

$\sigma^{*}=0.50000000017500$ If we choose level $\delta=0.8$ then
$\mathrm{x}_{* *}=(0.32,0.68)$
$\mu^{\mathrm{M}}\left(\mathrm{v}_{* *}^{\mathrm{M}}\right)=0.8, \mu^{\mathrm{L}}\left(\mathrm{v}_{* *}^{\mathrm{L}}\right)=0.4666667, \mu^{\mathrm{U}}\left(\mathrm{v}_{* *}^{\mathrm{U}}\right)=0.46666667$.

## Multi-level approach

## TFN's <br> (Elements of the pay-off matrix are triangular fuzzy numbers)

## BLLP

(Bi-level linear programming)

## MLLP

## FN's

(Elements of the pay-off matrix are fuzzy numbers)
(Multi-level linear programming)

## Multi-level linear programming



Thank you for your attention!

